

Classical Planning: Limits

Instantaneous actions

No temporal constraints

No concurrent actions

No continuous quantities

Spacecraft Domain

Observation-1
priority
time window
target
instruments
duration

Observation-2

Observation-3

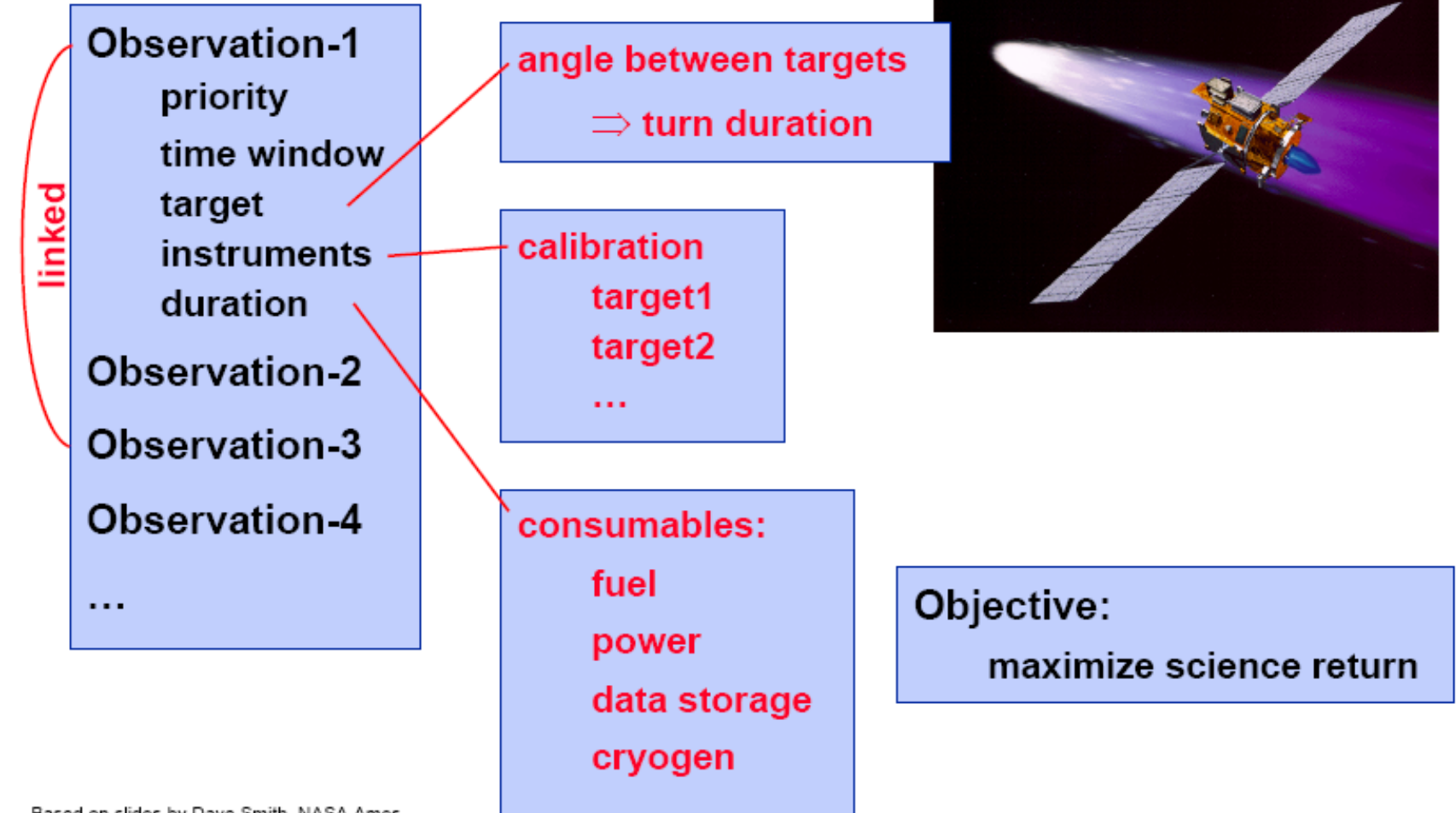
Observation-4

...



Objective:
maximize science return

Spacecraft Domain



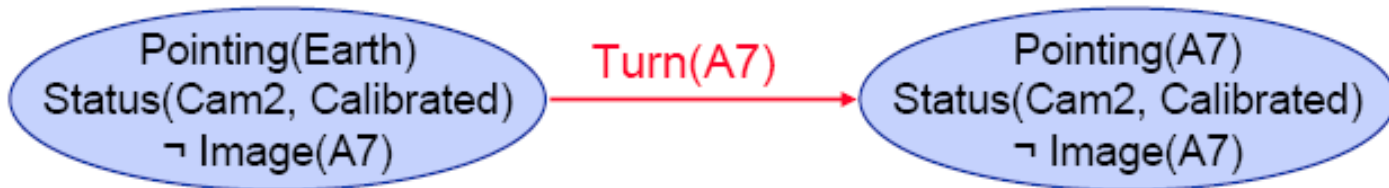
Extensions

- Time
- Resources
- Constraints
- Uncertainty
- Utility
- ...

Model

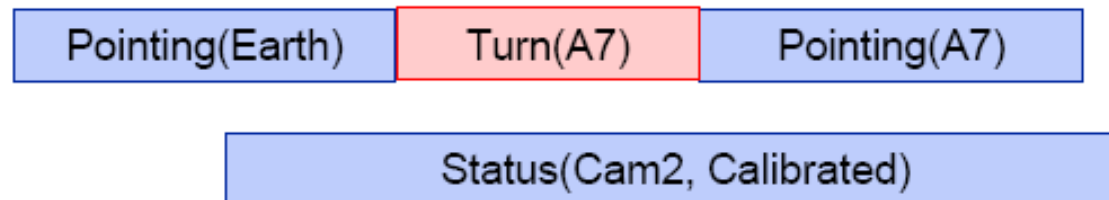
State-centric (McCarthy):

for each time describe propositions that are true



History-based (Hayes):

for each proposition describe times it is true



Temporal Interval Relations

A before B



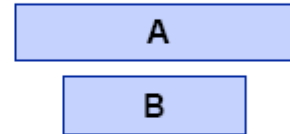
A meets B



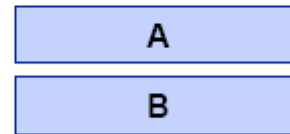
A overlaps B



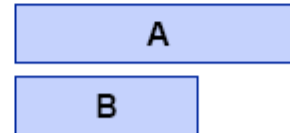
A contains B



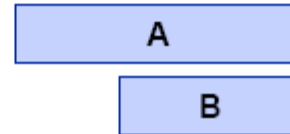
A = B



A starts B


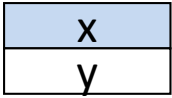

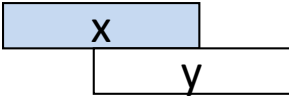
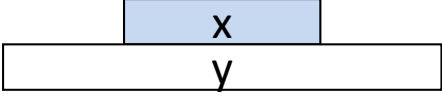
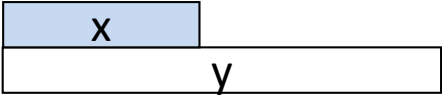



A ends B



Interval Algebra (aka Allen Algebra)

[Allen 83]

Relation	Symbol	Inverse	Illustration
X before Y	b	bi	
X equal Y	=		
X meets Y	m	mi	
X overlaps Y	o	oi	
X during Y	d	di	
X starts Y	s	si	
X finishes Y	f	fi	

Interval Algebra: Qualitative TN

- Variables
 - An interval represent an event with some duration
- Constraints
 - Intervals I, J are related by a binary constraint
 - The constraint is a subset of the 13 basic relations
 $r = \{ b, m, o, s, d, f, bi, mi, oi, si, di, fi, = \}$
 - Example: $I \{r_1, r_2, \dots, r_k\} J \Leftrightarrow (I r_1 J) \vee (I r_2 J) \vee \dots \vee (I r_k J)$
 - Enumerate atomic relations between two variables

Interval Algebra Constraint Network

- Variables: temporal intervals I and J
- Domain: set of ordered pairs of real numbers
- Constraints are subsets of the 13 relations
 - How many distinct relations?
- A solution is an assignment of a pair of numbers to each variable such that no constraint is violated

Interval Algebra: Example

Story:

John was not in the room when I touched the switch to turn on the light but John was in the room later when the light was on.

CSP model:

Variables:

Switch – the time of touching the switch

Light – the light was on

Room – the time that John was in the room

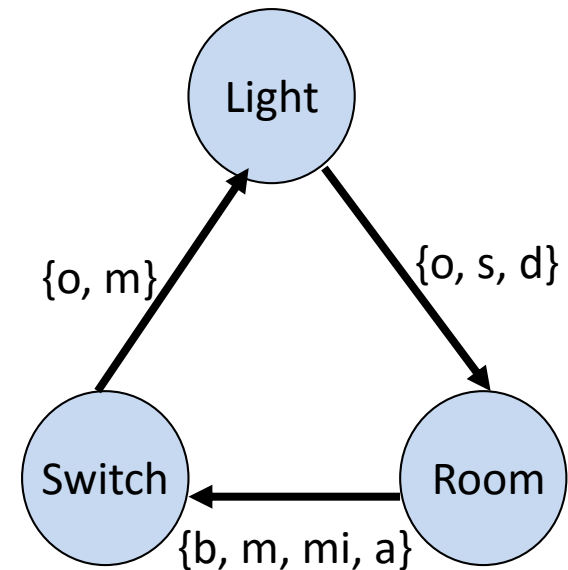
Constraints:

Switch overlaps or meets Light: $S \{o, m\} L$

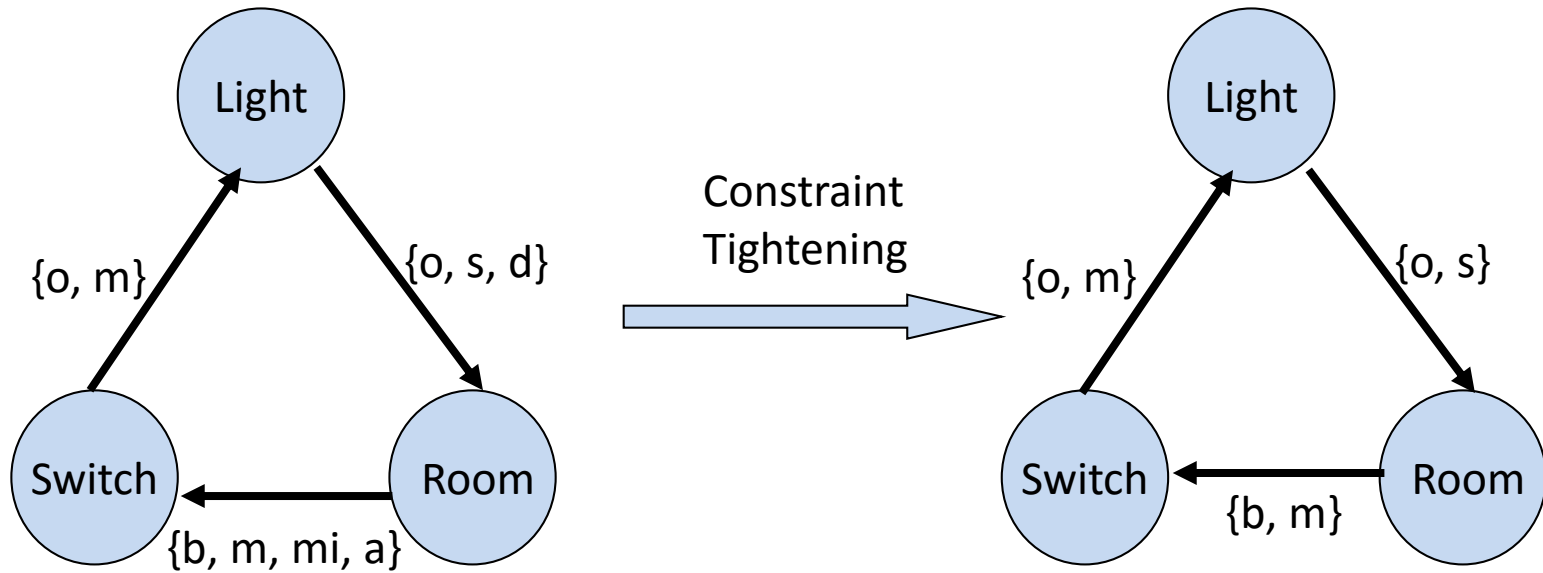
Switch is before, meets, is met by or after

Room: $S \{b, m, mi, bi\} R$

Light overlaps, starts or is during Room: $L \{o, s, d\} R$



The Task: Get the Minimal Network



- A unique network equivalent to original network
- All constraints are subsets of original constraints
- Provides a more explicit representation
- Useful in answering many types of queries

Temporal Operators

TakeImage (?target, ?instr):

Pre: Status(?instr, Calibrated), Pointing(?target)

Eff: Image(?target)



TakeImage (?target, ?instr)

contained-by

Status(?instr, Calibrated)

contained-by

Pointing(?target)

meets

Image(?target)

Temporal Operators

TakelImage (?target, ?instr)

contained-by

contained-by

meets

Status(?instr, Calibrated)

Pointing(?target)

Image(?target)



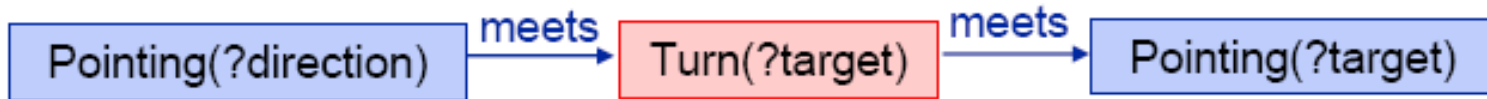
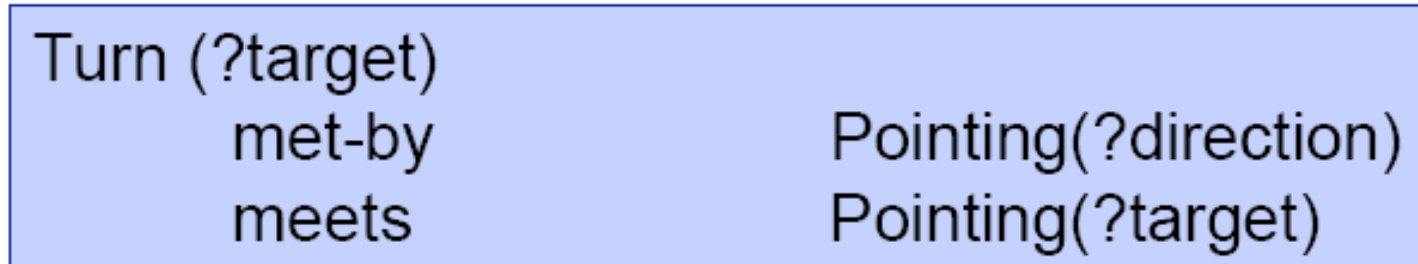
Temporal Operators

TakelImage (?target, ?instr)	
contained-by	Status(?instr, Calibrated)
contained-by	Pointing(?target)
meets	Image(?target)



$\text{TakelImage}(\text{?target}, \text{?instr})_A$
 $\Rightarrow \exists P \{ \text{Status}(\text{?instr}, \text{Calibrated})_P \wedge \text{Contains}(P, A) \}$
 $\wedge \exists Q \{ \text{Pointing}(\text{?target})_Q \wedge \text{Contains}(Q, A) \}$
 $\wedge \exists R \{ \text{Image}(\text{?target})_R \wedge \text{Meets}(A, R) \}$

Temporal Operators



Temporal Operators

Calibrate (?instr)

met-by

contained-by

contained-by

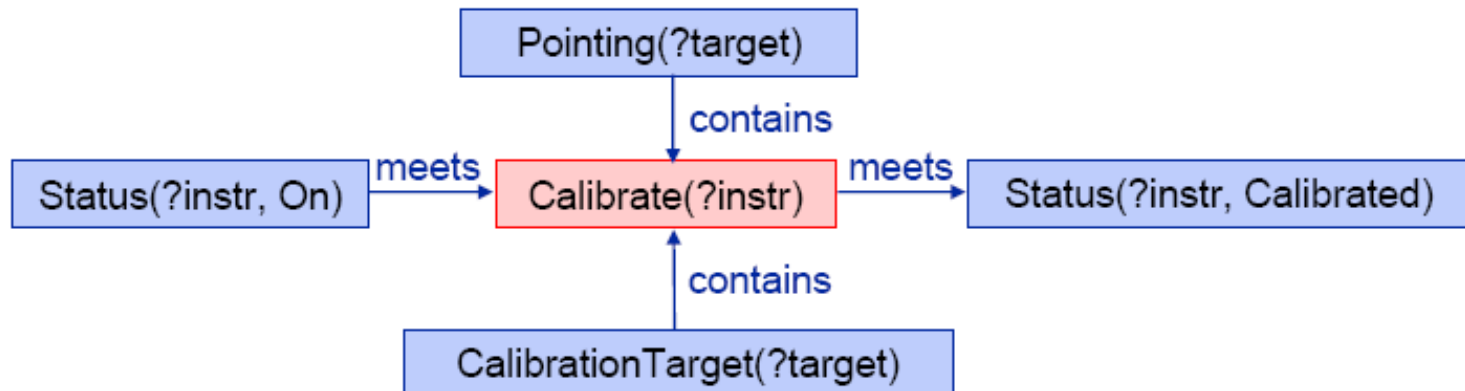
meets

Status(?instr, On)

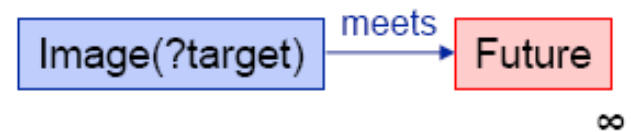
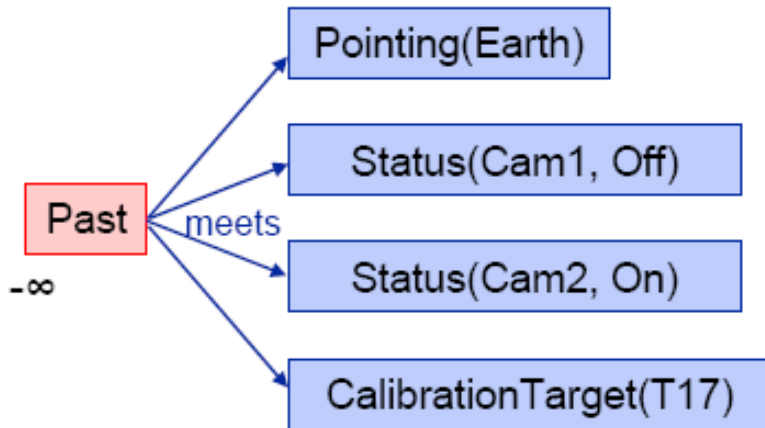
CalibrationTarget(?target)

Pointing(?target)

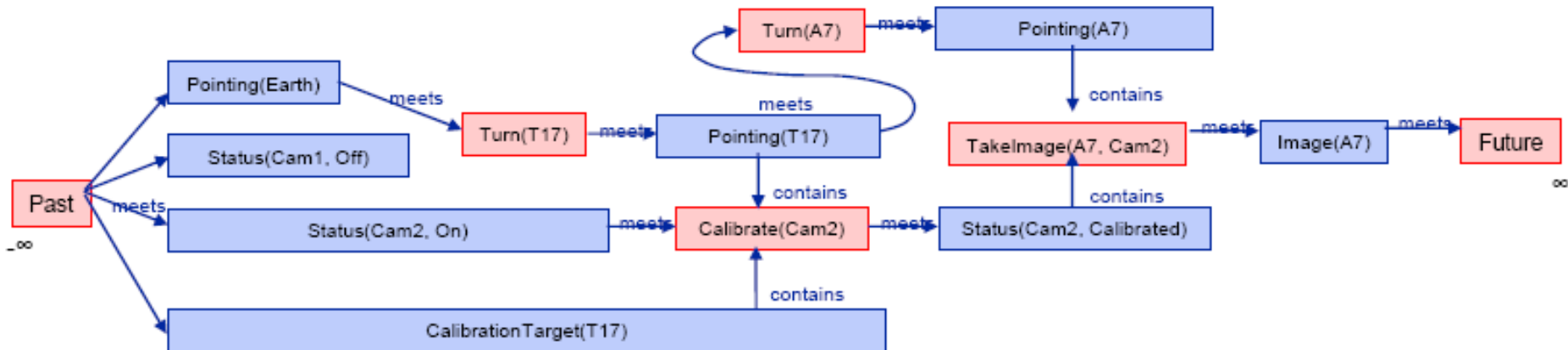
Status(?instr, Calibrated)



Temporal Planning Problem



Consistent Complete Plan



CBI-Planning

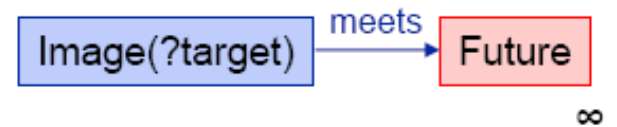
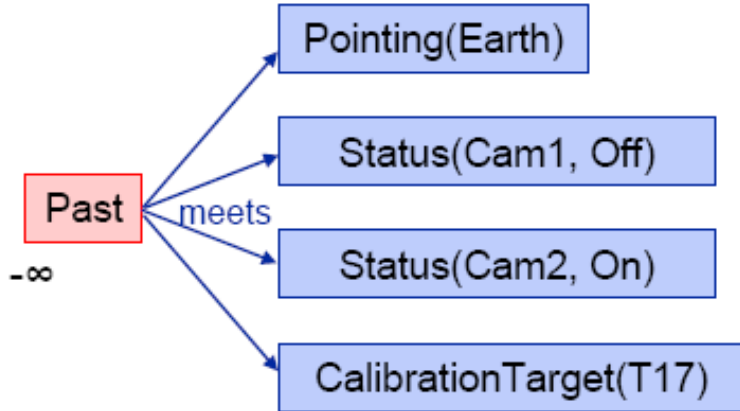
Choose:

- introduce an action & instantiate constraints

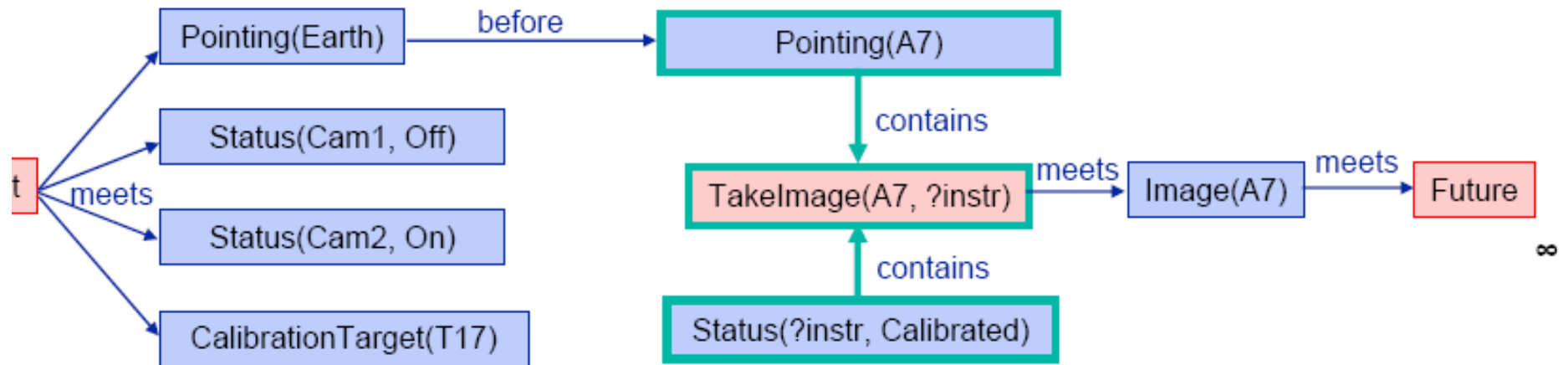
- coalesce propositions

Propagate constraints

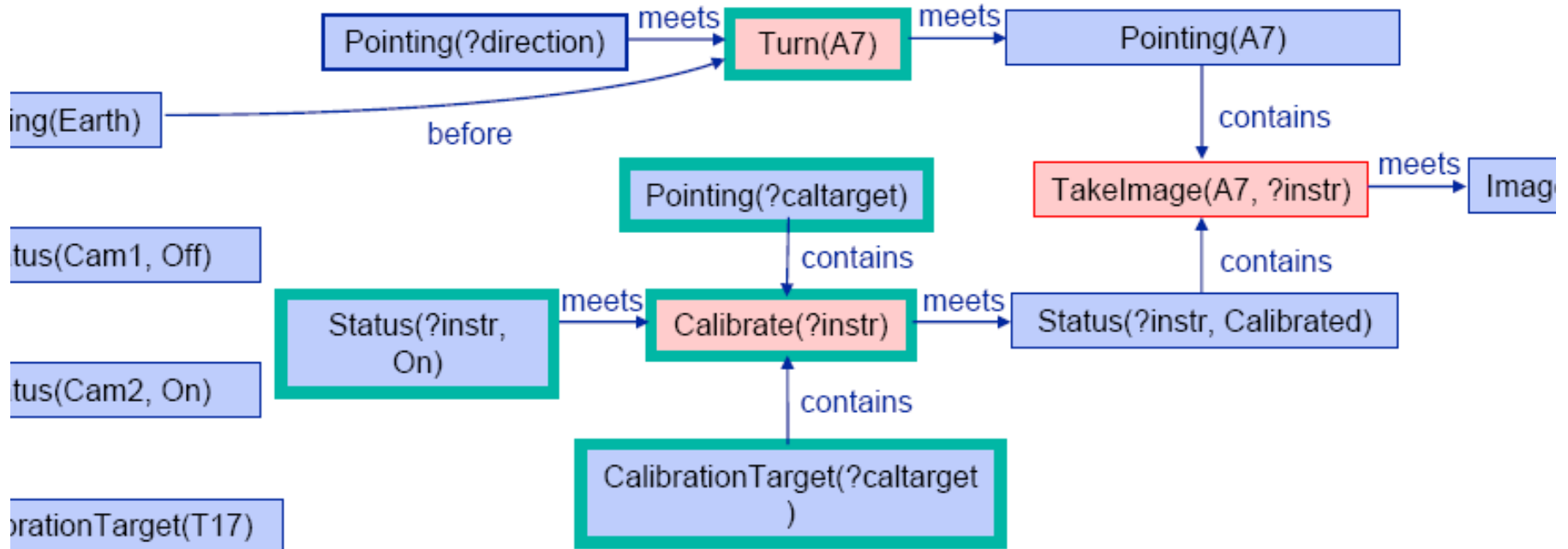
Initial Plan



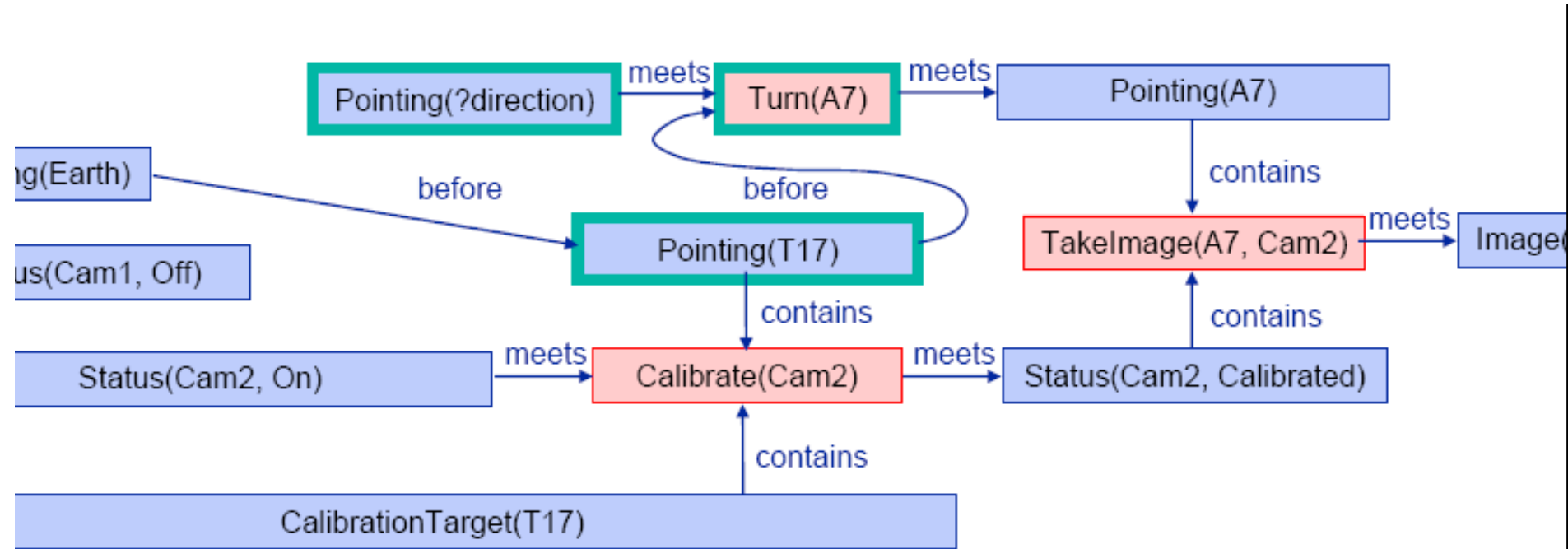
Expansion



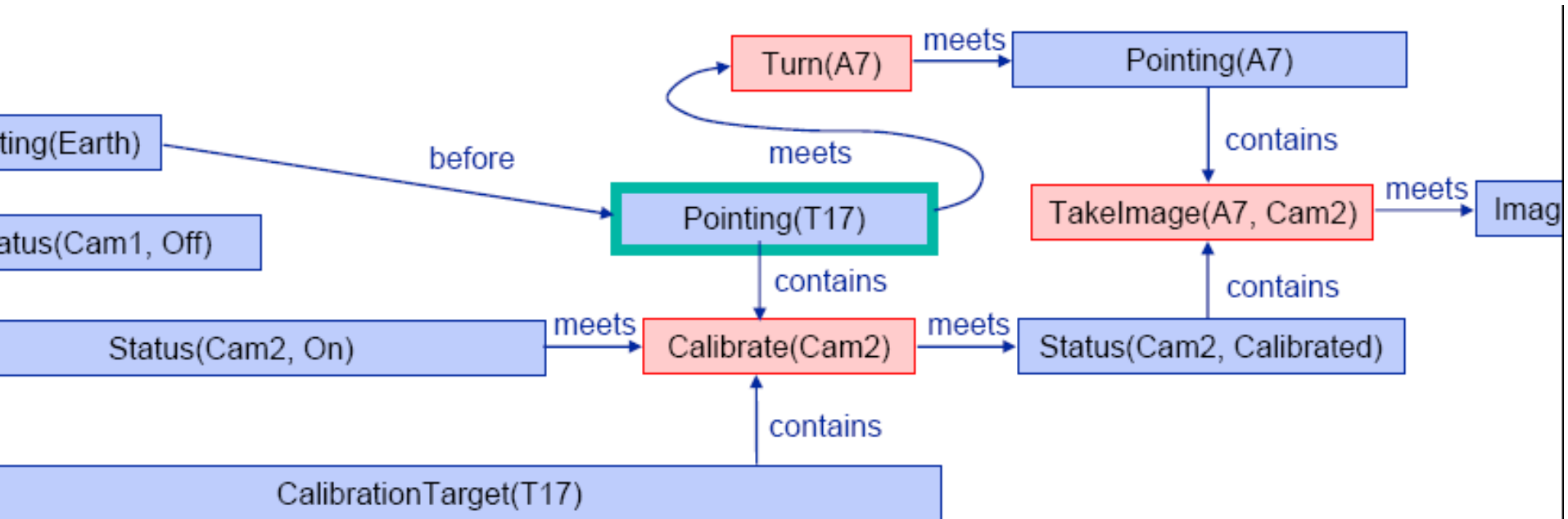
Expansion



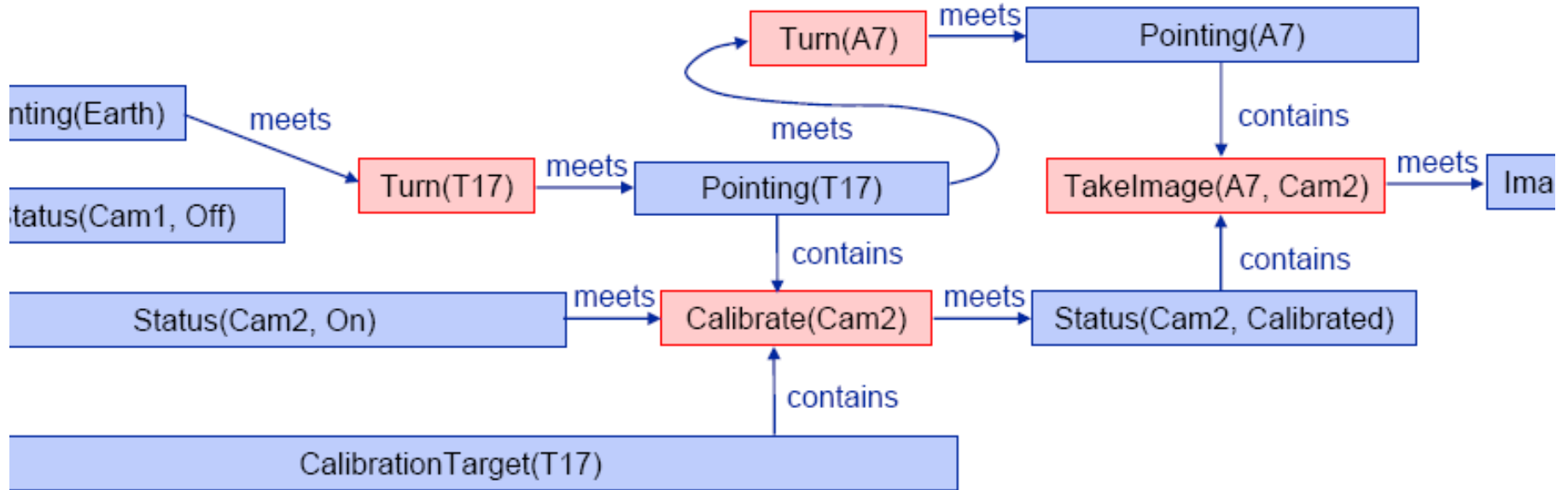
Coalescing



Coalescing



Coalescing



CBI-Algorithm

Expand(TQAs, constraints)

1. If the constraints are inconsistent, **fail**
2. If all TQAs have causal explanations, **return**(TQAs, constraints)
3. Select a $g \in \text{TQAs}$ with no causal explanation
4. **Choose**:
 - Choose** another $p \in \text{TQAs}$ such that g can be coalesced with p under constraints C
Expand(TQAs- g , constraints $\cup C$)
 - Choose** an action that would provide a causal explanation for g
Let A be a new TQA for the action,
and let R be the set of new TQAs implied by the axioms for A
Let C be the constraints between A and R
Expand(TQAs $\cup \{A\} \cup R$, constraints $\cup C$)

CBI-Planners

Zeno (Penberthy)

intervals, no CSP

Trains (Allen)

Descartes (Joslin)

extreme least commitment

IxTeT (Ghallab)

functional rep.

HSTS (Muscettola)

functional rep., activities

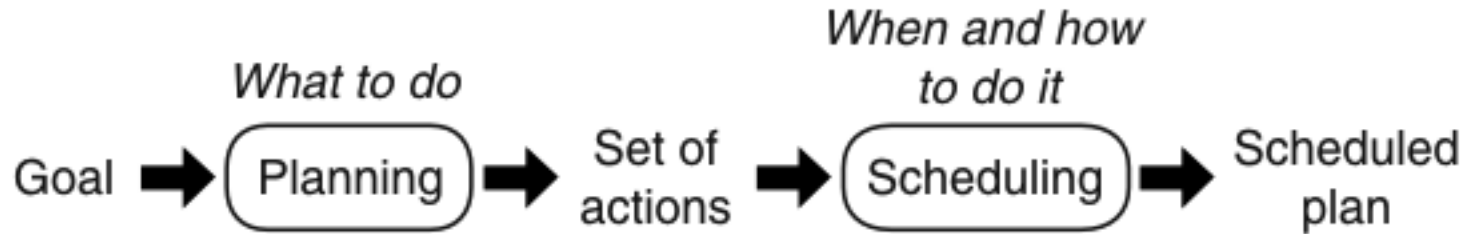
EUROPA (Jonsson)

functional rep., activities

CBI vs POP

- CBI is similar to POP because least commitment and partial order
- But, temporal constraints in CBI ...
- Constraints Temporal Network associated with a plan
- Constraint propagation

Planning and Scheduling

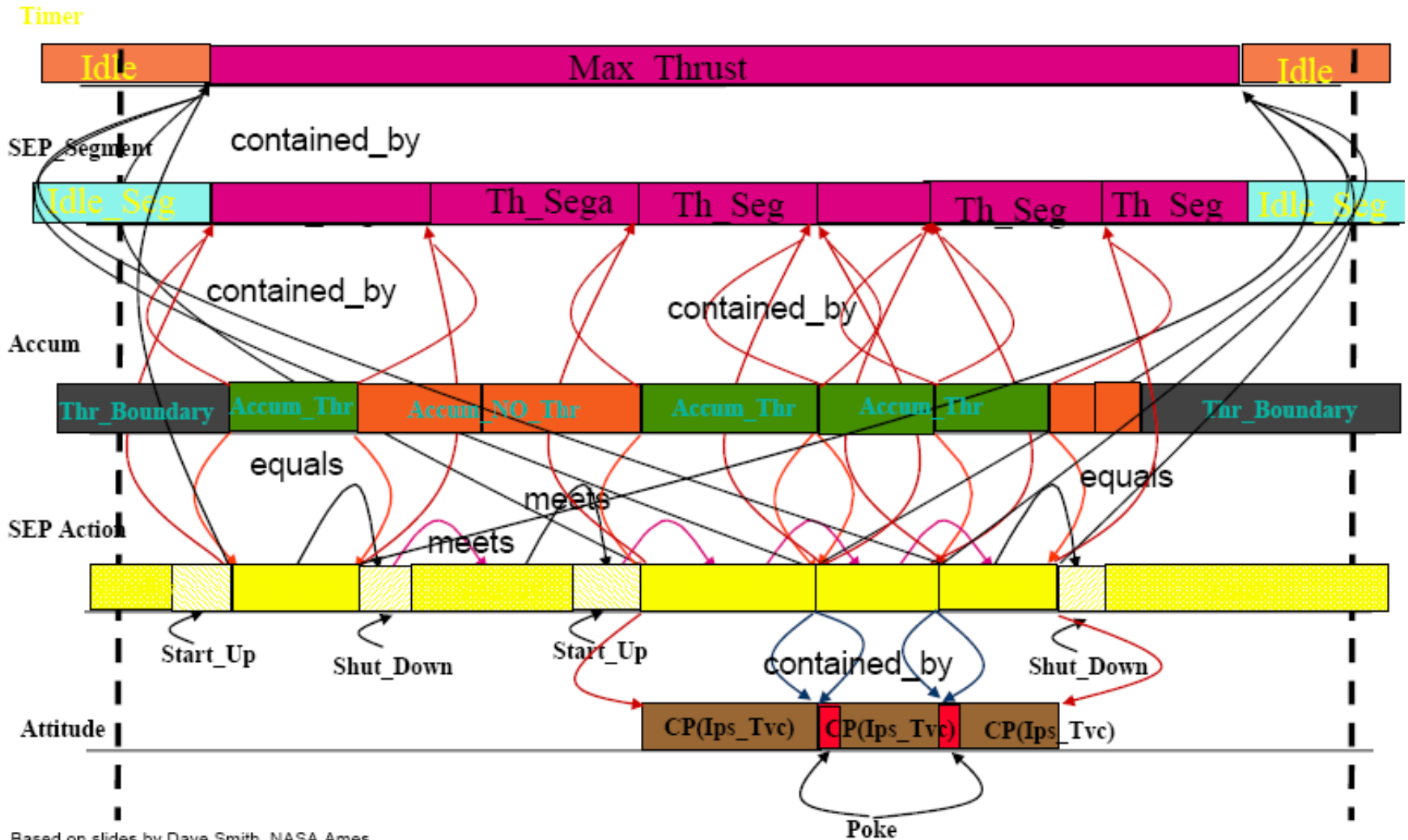


- Scheduling has usually been addressed separately from planning
- Thus, will give an overview of scheduling algorithms
- In some cases, cannot decompose planning and scheduling so cleanly

Temporal Constraints

- x before y
 - x meets y
 - x overlaps y
 - x during y
 - x starts y
 - x finishes y
 - x equals y
-
- y after x
 - y met-by x
 - y overlapped-by x
 - y contains x
 - y started-by x
 - y finished-by x
 - y equals x

RAX Example: DS1



Temporal Constraints as Inequalities

- x before y $X^+ < Y^-$
- x meets y $X^+ = Y^-$
- x overlaps y $(Y^- < X^+) \ \& \ (X^- < Y^+)$
- x during y $(Y^- < X^-) \ \& \ (X^+ < Y^+)$
- x starts y $(X^- = Y^-) \ \& \ (X^+ < Y^+)$
- x finishes y $(X^- < Y^-) \ \& \ (X^+ = Y^+)$
- x equals y $(X^- = Y^-) \ \& \ (X^+ = Y^+)$

Inequalities may be expressed as binary interval relations:

$$X^+ - Y^- < [-\text{inf}, 0]$$

Metric Constraints

- Going to the store takes at least 10 minutes and at most 30 minutes.
→ $10 \leq [T^+(\text{store}) - T^-(\text{store})] \leq 30$
- Bread should be eaten within a day of baking.
→ $0 \leq [T^+(\text{baking}) - T^-(\text{eating})] \leq 1 \text{ day}$
- Inequalities, $X^+ < Y^-$, may be expressed as binary interval relations:
→ $-\text{inf} < [X^+ - Y^-] < 0$

Temporal Constraint Networks

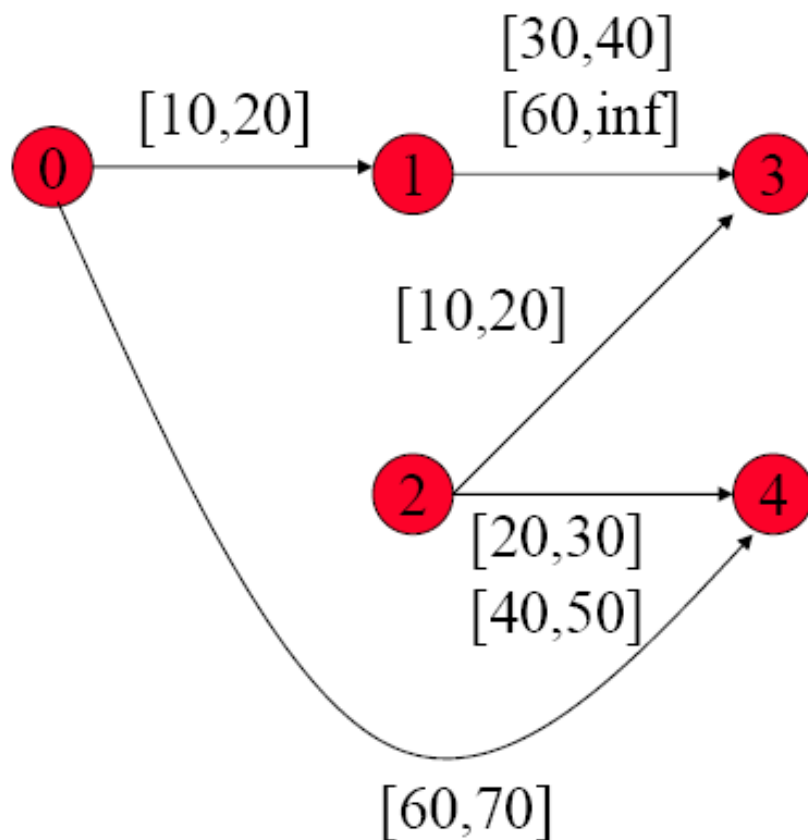
- A set of time points X_i at which events occur.
- Unary constraints

$$(a_0 \leq X_i \leq b_0) \text{ or } (a_1 \leq X_i \leq b_1) \text{ or } \dots$$

- Binary constraints

$$(a_0 \leq X_j - X_i \leq b_0) \text{ or } (a_1 \leq X_j - X_i \leq b_1) \text{ or } \dots$$

Temporal Constraint Satisfaction Problem



Simple Temporal Networks

Simple Temporal Networks:

- A set of time points X_i at which events occur.

- Unary constraints

$$(a_0 \leq X_i \leq b_0) \text{ or } (a_1 \leq X_i \leq b_1) \text{ or } \dots$$

- Binary constraints

$$(a_0 \leq X_j - X_i \leq b_0) \text{ or } (a_1 \leq X_j - X_i \leq b_1) \text{ or } \dots$$

Sufficient to represent:

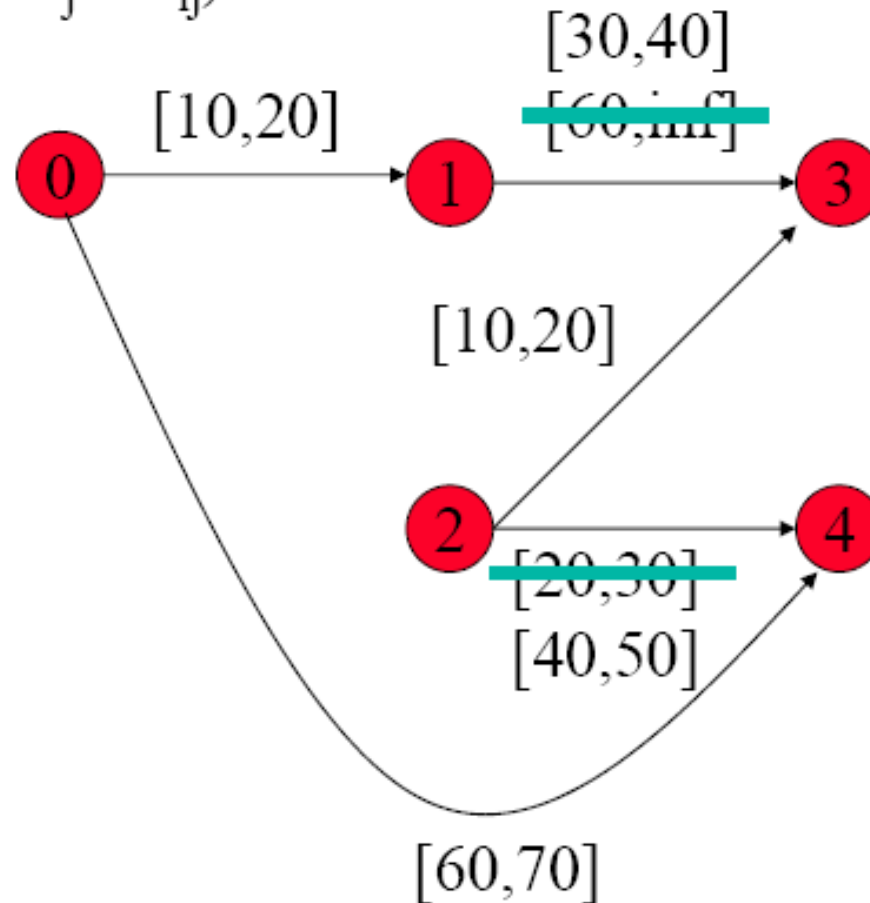
- most Allen relations
- simple metric constraints

Can't represent:

- Disjoint activities

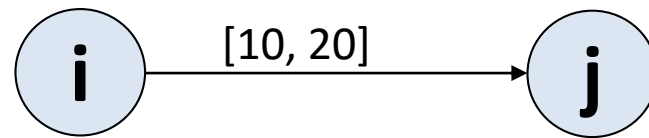
Simple Temporal Networks

- $T_{ij} = (a_{ij} \leq X_i - X_j \leq b_{ij})$



Simple Temporal Network (STP)

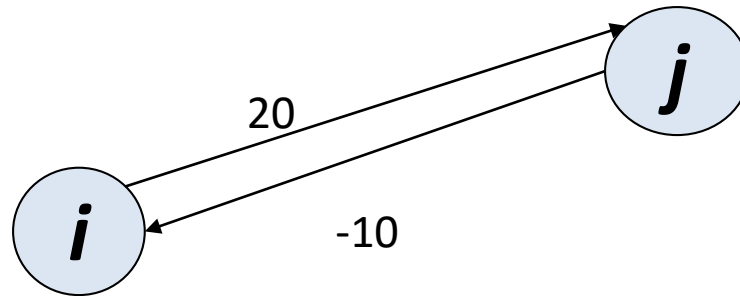
- A special class of temporal problems
- Can be solved in polynomial time
- An edge $e_{ij}: i \rightarrow j$ is labeled by a **single** interval $[a_{ij}, b_{ij}]$



- Constraint $(a_{ij} \leq x_j - x_i \leq b_{ij})$ expressed by
 $(x_j - x_i \leq b_{ij}) \wedge (x_i - x_j \leq -a_{ij})$
- Example $(x_j - x_i \leq 20) \wedge (x_i - x_j \leq -10)$

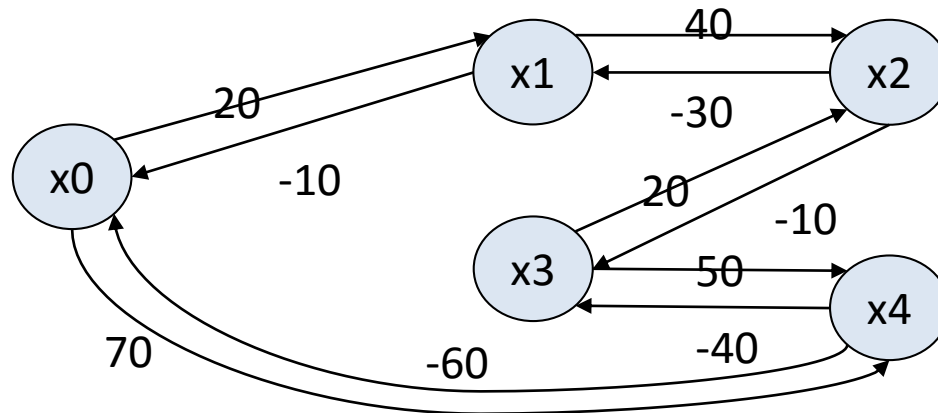
Distance Graph of an STP

- The STP is transformed into an all-pairs-shortest-paths problem on a **distance graph**
- Each constraint is replaced by two edges: one + and one -



- Constraint graph \rightarrow directed cyclic graph

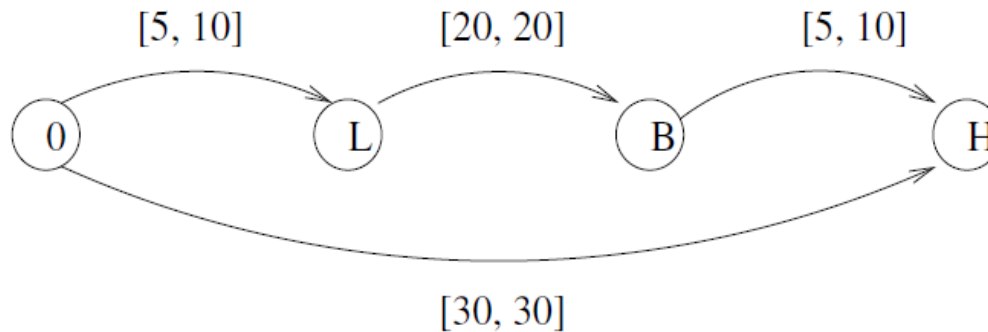
Solving the Distance Graph of the STP



- Run **Floyd-Warshall** all pairs shortest path
- If any pair of nodes has a negative cycle \Rightarrow inconsistency
- If consistent after **F-W** \Rightarrow minimal & decomposable
- Once d-graph formed, assembling a solution by checking against the previous labeling
- Total time: F-W $O(n^3)$ + Assembling $O(n^2)$ = $O(n^3)$.

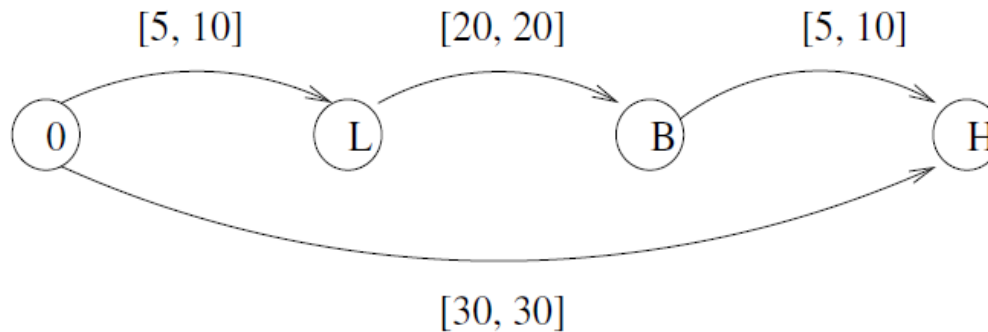
Example

- **Eventi:**
 1. I was in Houghton at 8:30.
 2. I left home between 8:05 and 8:10.
 3. It takes me 20 minutes to drive to the bridge.
 4. I waited 5-10 minutes at the bridge.



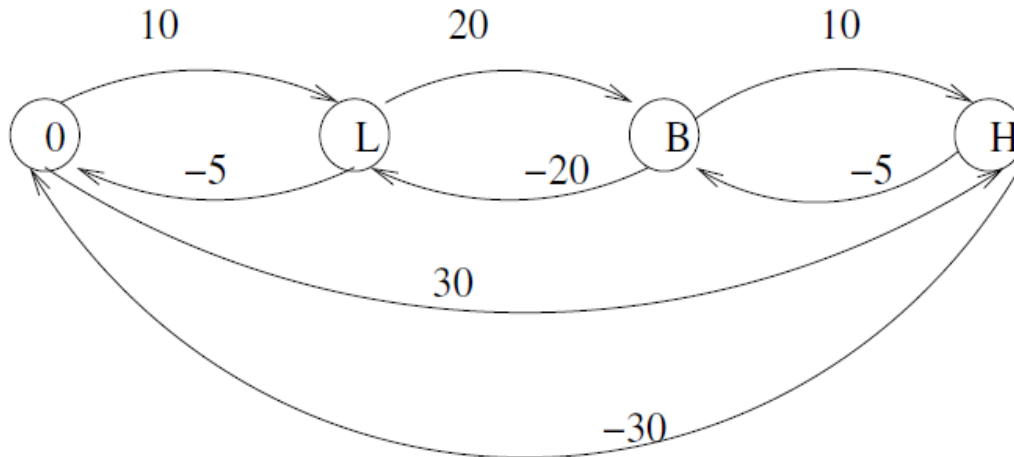
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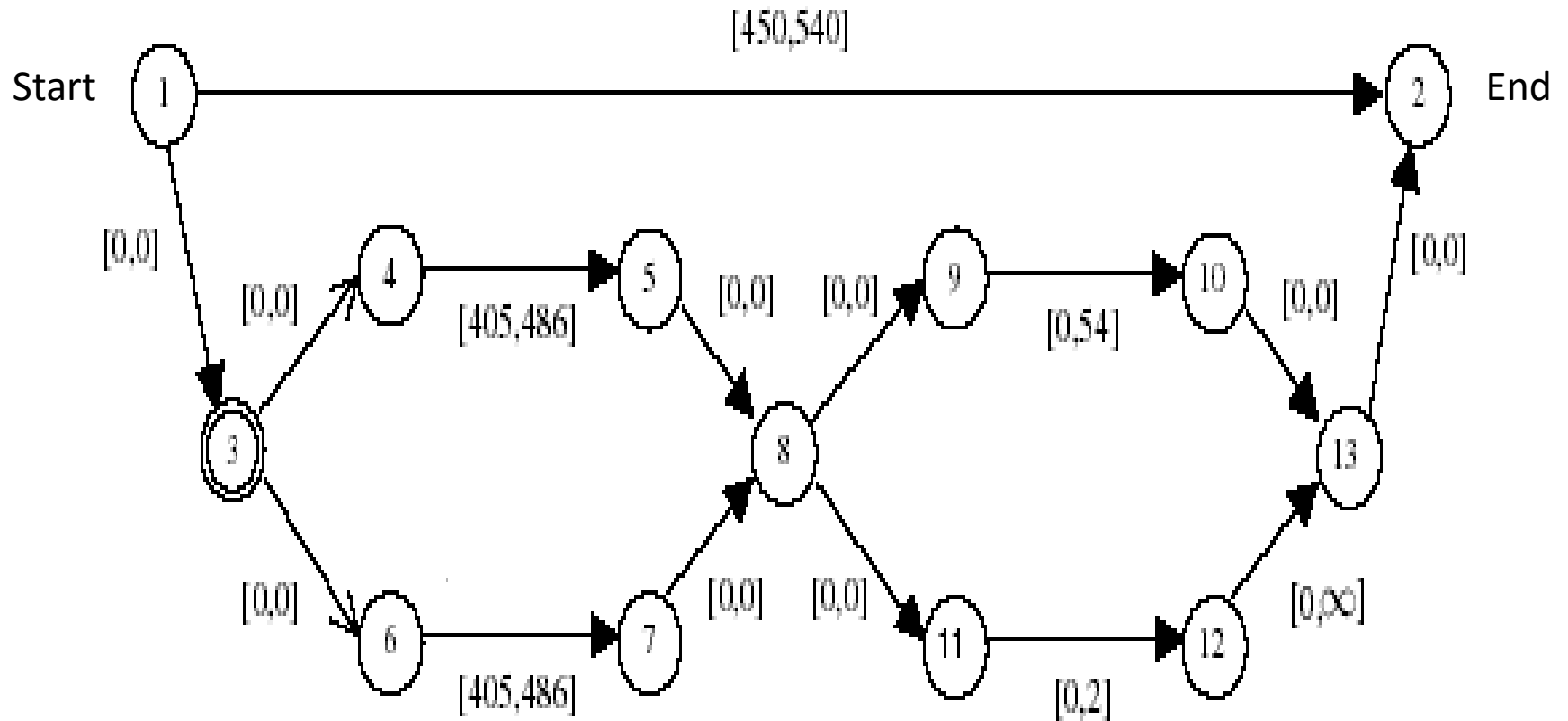


	0	1	2	3
0	0	10	99	30
1	-5	0	20	99
2	99	-20	0	10
3	-30	99	-5	0

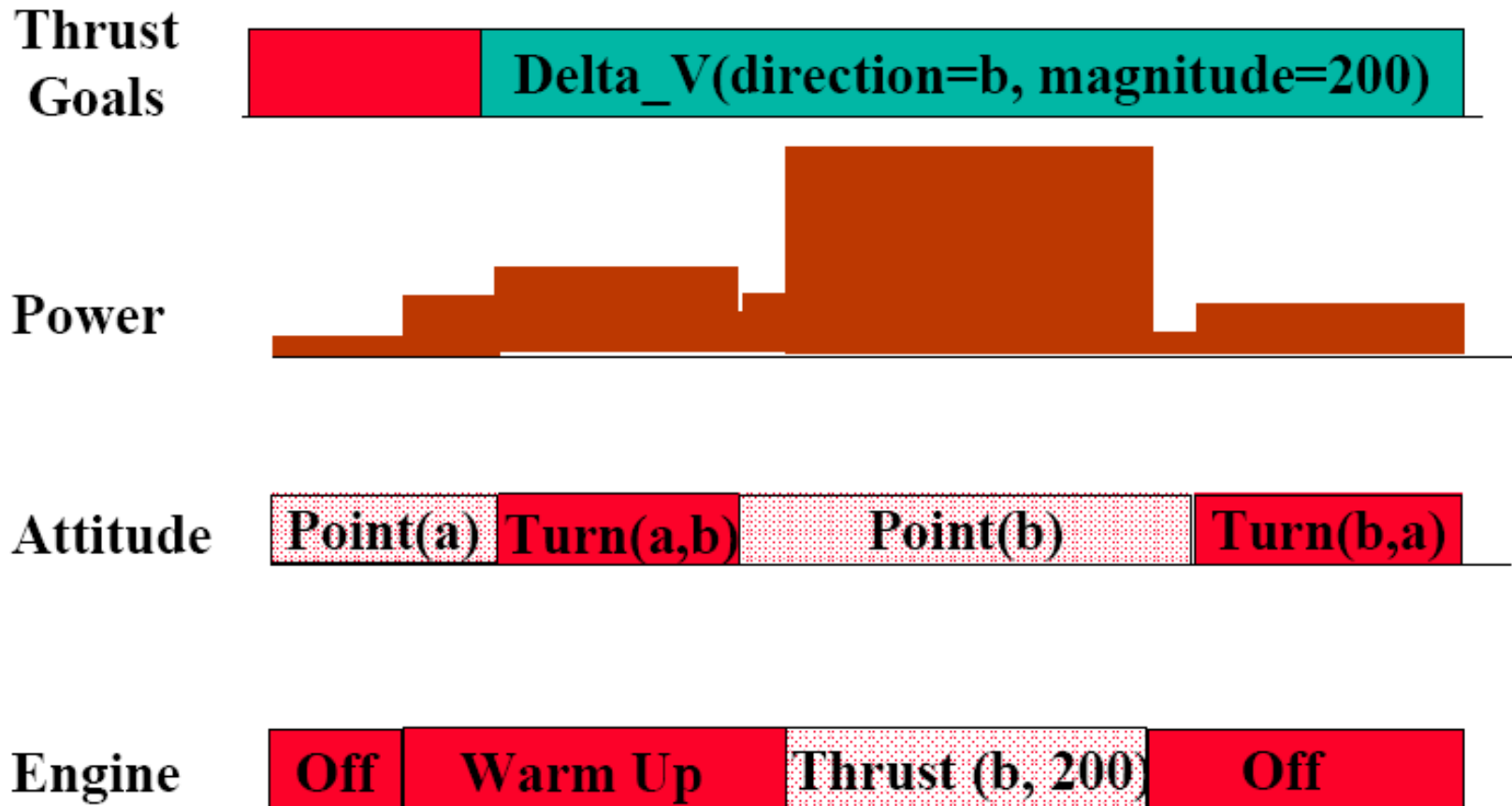
Floyd-Warshall

	0	1	2	3
0	0	5	25	30
1	-5	0	20	25
2	-25	-20	0	5
3	-30	-25	-5	0

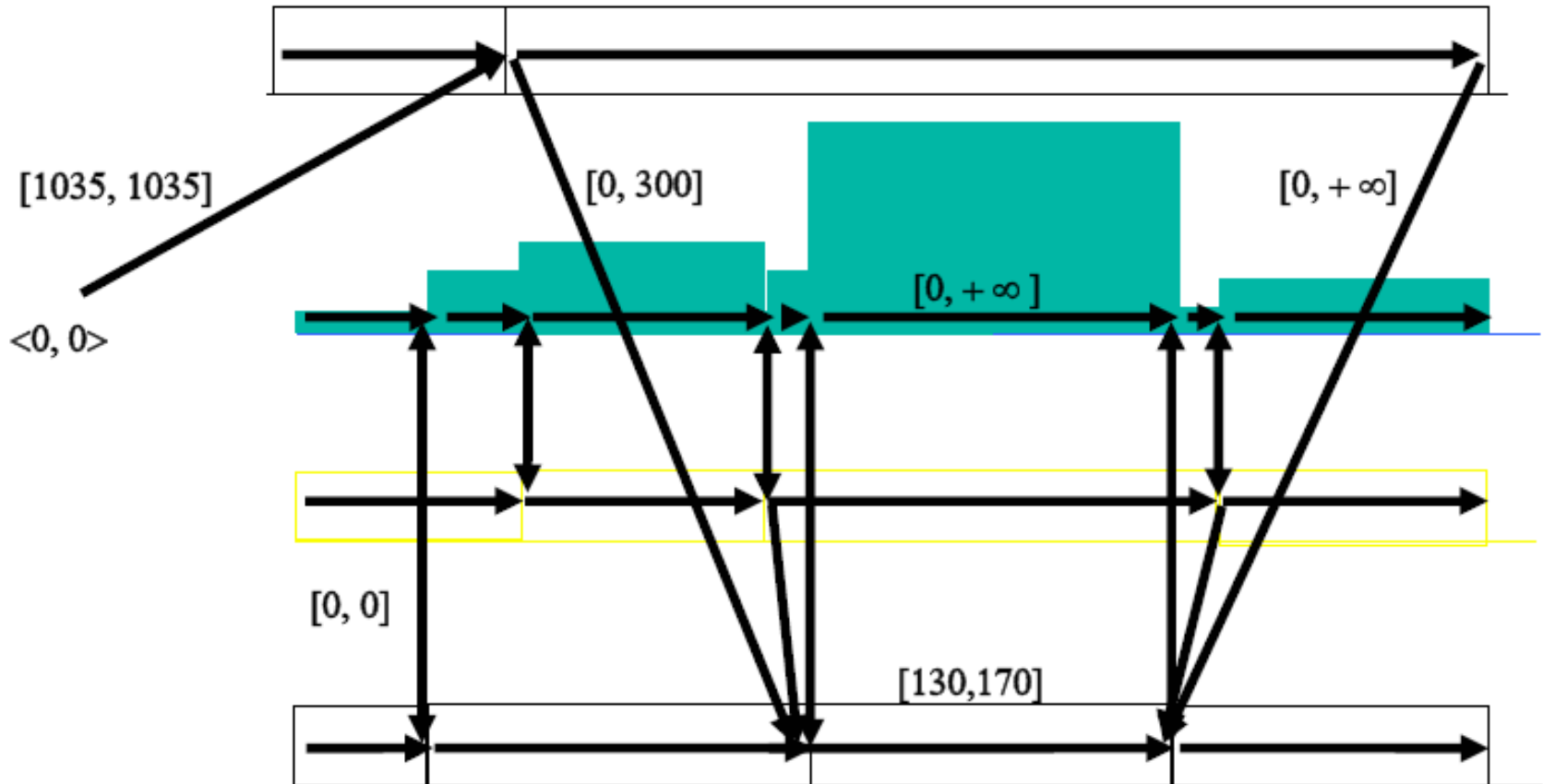
STN example



A Complete CBI-Plan is a STN



A Complete CBI-Plan is a STN



DS1: Remote Agent

Remote Agent on Deep Space 1



Started: January 1996
Launch: Fall 1998

Remote Agent Experiment: RAX

Remote Agent Experiment

See rax.arc.nasa.gov

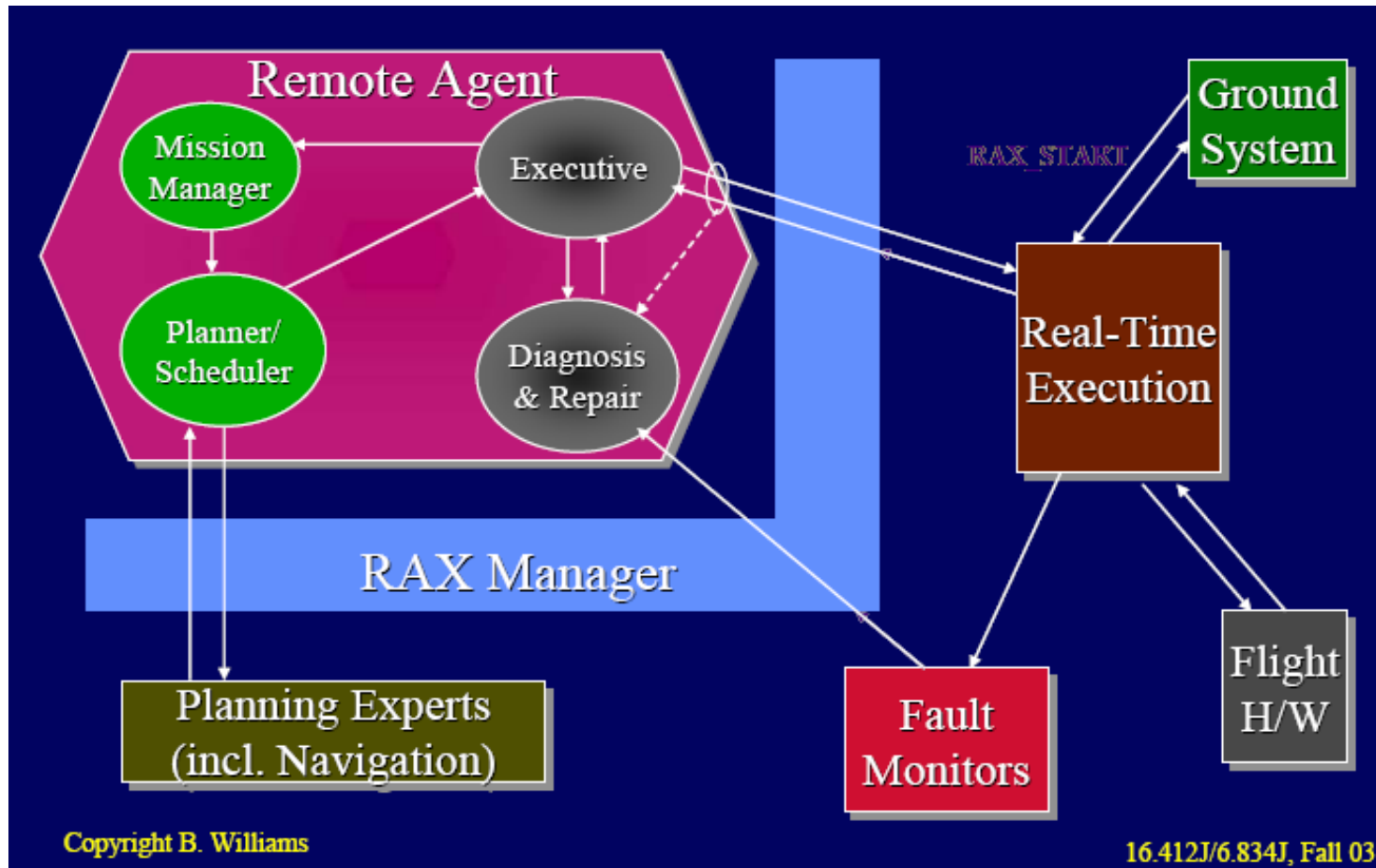
May 17-18th experiment

- Generate plan for course correction and thrust
- **Diagnose camera as stuck on**
 - **Power constraints violated, abort current plan and replan**
- Perform optical navigation
- Perform ion propulsion thrust

May 21th experiment.

- **Diagnose faulty device and**
 - **Repair by issuing reset.**
- **Diagnose switch sensor failure.**
 - **Determine harmless, and continue plan.**
- **Diagnose thruster stuck closed and**
 - **Repair by switching to alternate method of thrusting.**
- Back to back planning

Remote Agent



Remote Agent

**Thrust
Goals**

Power

Attitude

Engine

Remote Agent

- Mission Manager

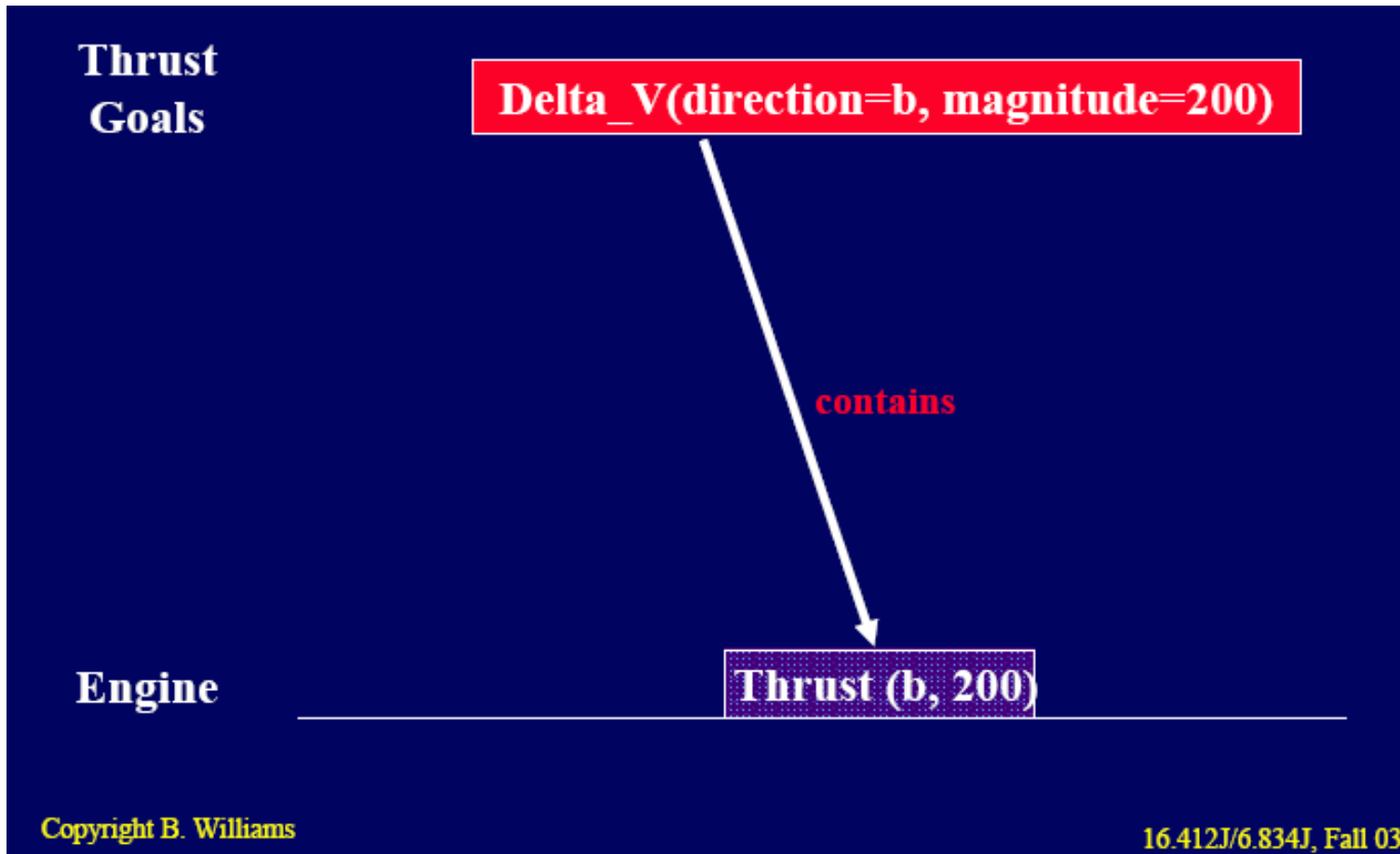
The image shows a Mission Manager interface with a dark blue background. It features four main sections: Thrust Goals, Power, Attitude, and Engine. Each section has a horizontal bar representing its status or value.

- Thrust Goals:** A horizontal bar with a blue segment on the left and a red segment on the right. The red segment contains the text "Delta_V(direction=b, magnitude=200)".
- Power:** A horizontal bar that is currently empty.
- Attitude:** A horizontal bar with a purple segment on the left containing the text "Point(a)".
- Engine:** A horizontal bar with two blue segments, each containing the text "Off".

Copyright B. Williams 16.412J/6.834J, Fall 03

Remote Agent

- Constraints:



Remote Agent

- Planner starts

The diagram illustrates the state of a Remote Agent across four categories, each with a horizontal bar representing its value or state:

- Thrust Goals:** A blue bar on the left and a red bar on the right containing the text `Delta_V(direction=b, magnitude=200)`.
- Power:** A single horizontal line with no text or markers.
- Attitude:** A purple dotted box containing the text `Point(a)` positioned at the start of a horizontal line.
- Engine:** Two blue boxes, each containing the text `Off`, positioned at the start and end of a horizontal line.

Copyright B. Williams 16.412J/6.834J, Fall 03

Remote Agent

- Planning

**Thrust
Goals**

Delta_V(direction=b, magnitude=200)

Power

Attitude

Point(a)

Engine

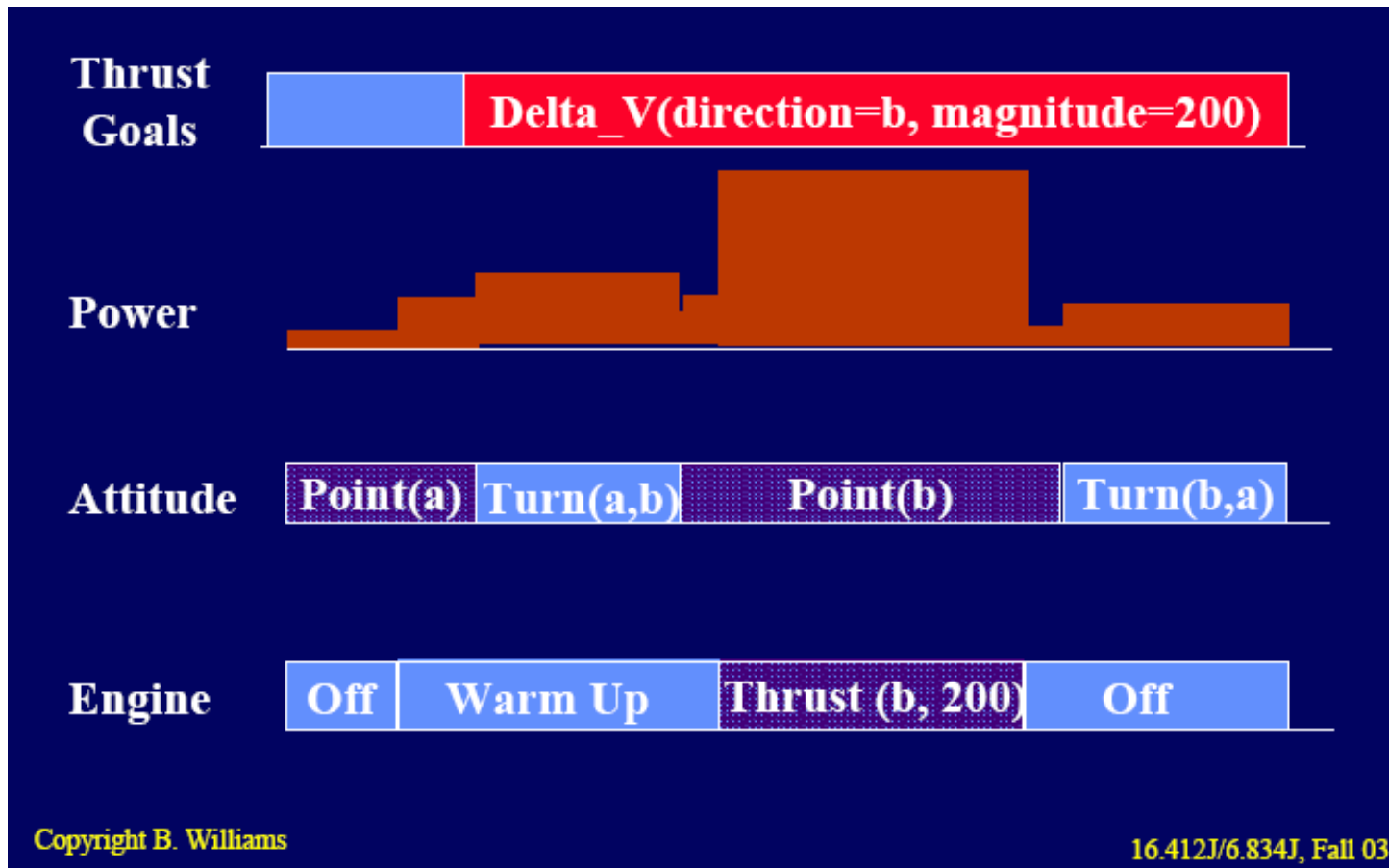
Off

Thrust (b, 200)

Off

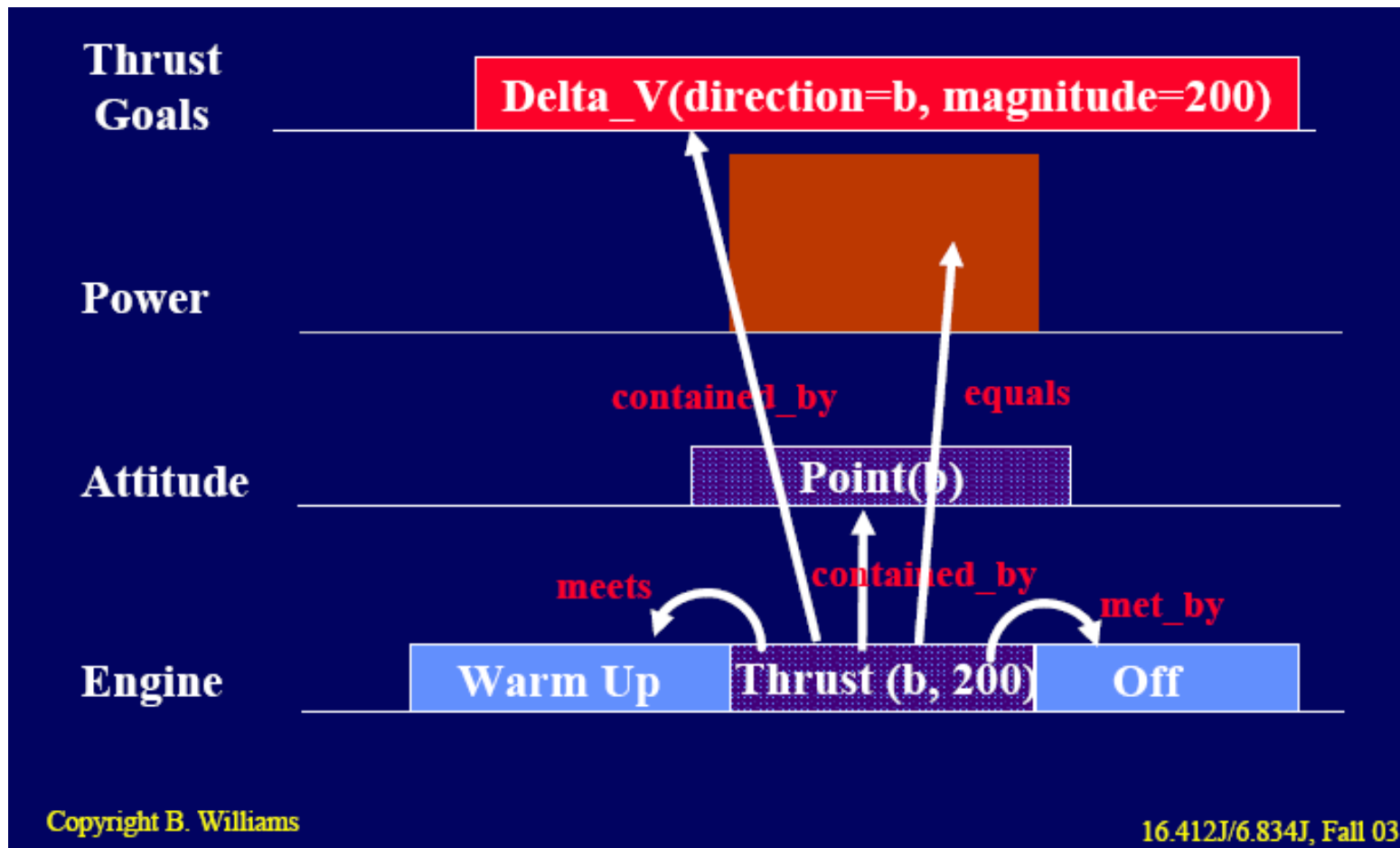
Remote Agent

- Final Plan



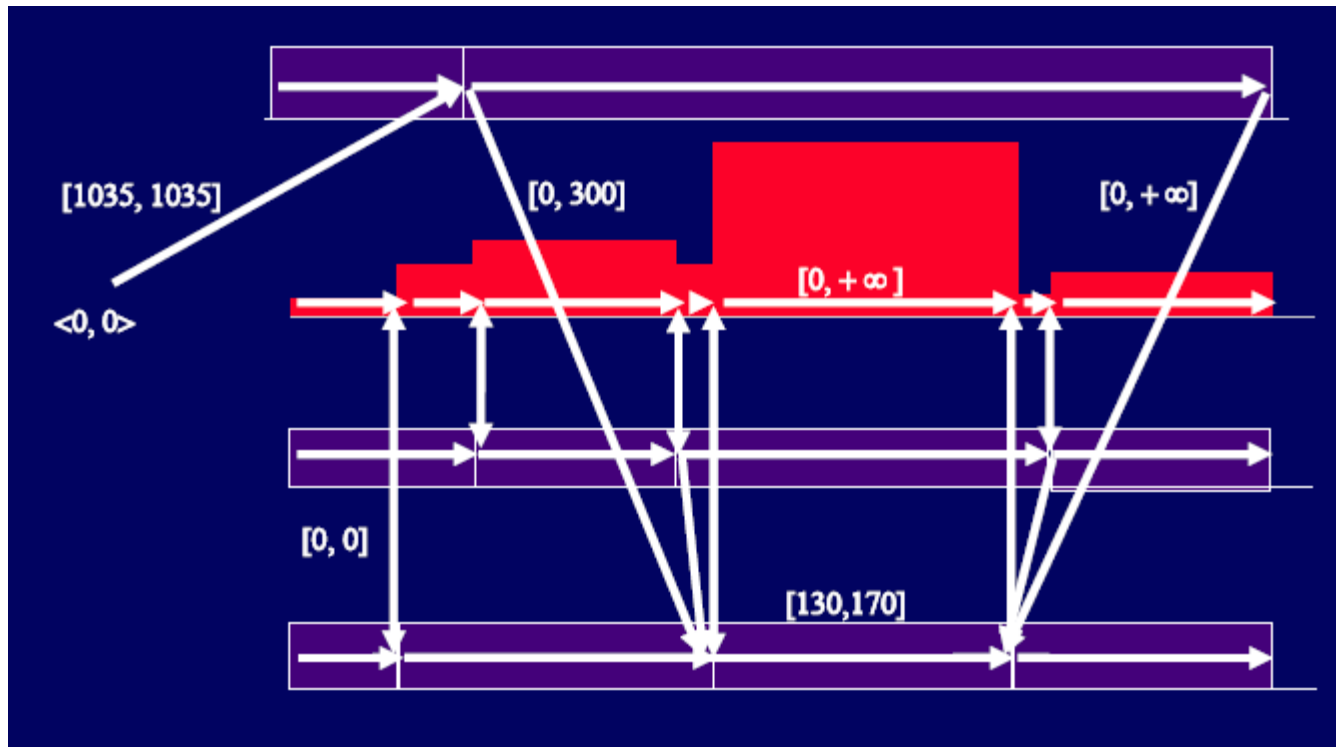
Remote Agent

- Constraints



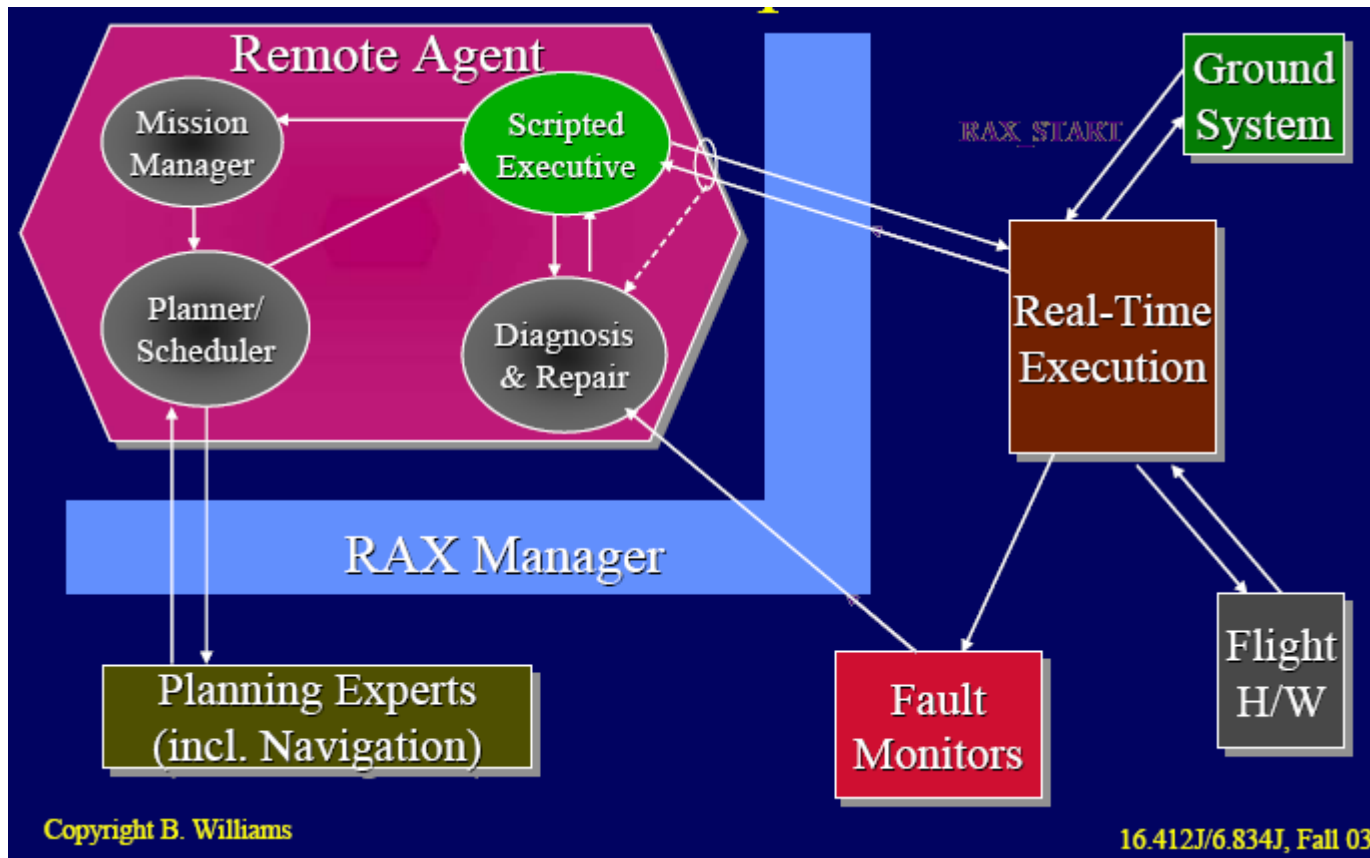
Remote Agent

- Flexible Temporal Plan through least commitment



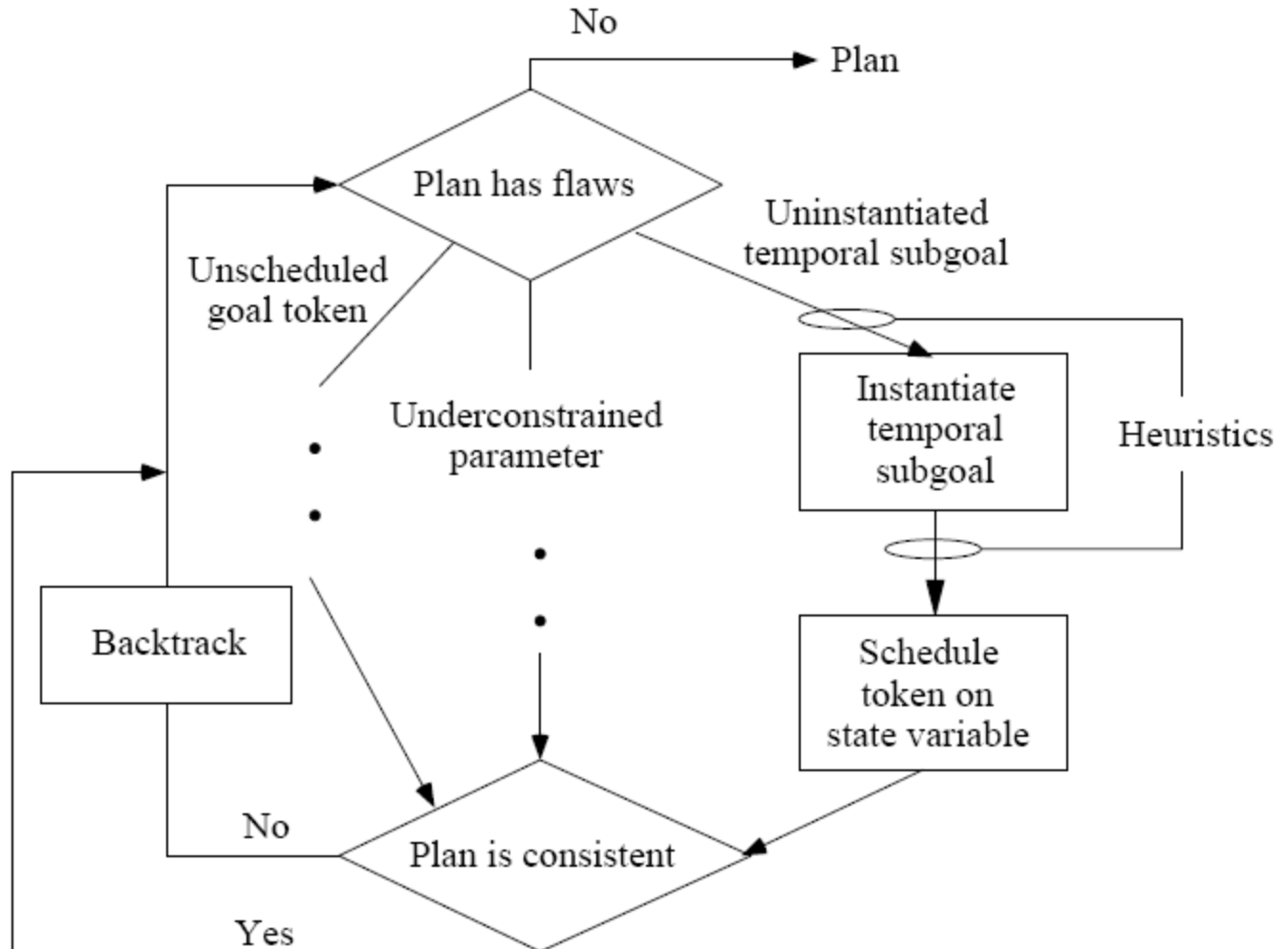
Remote Agent

- Executive system dispatch tasks



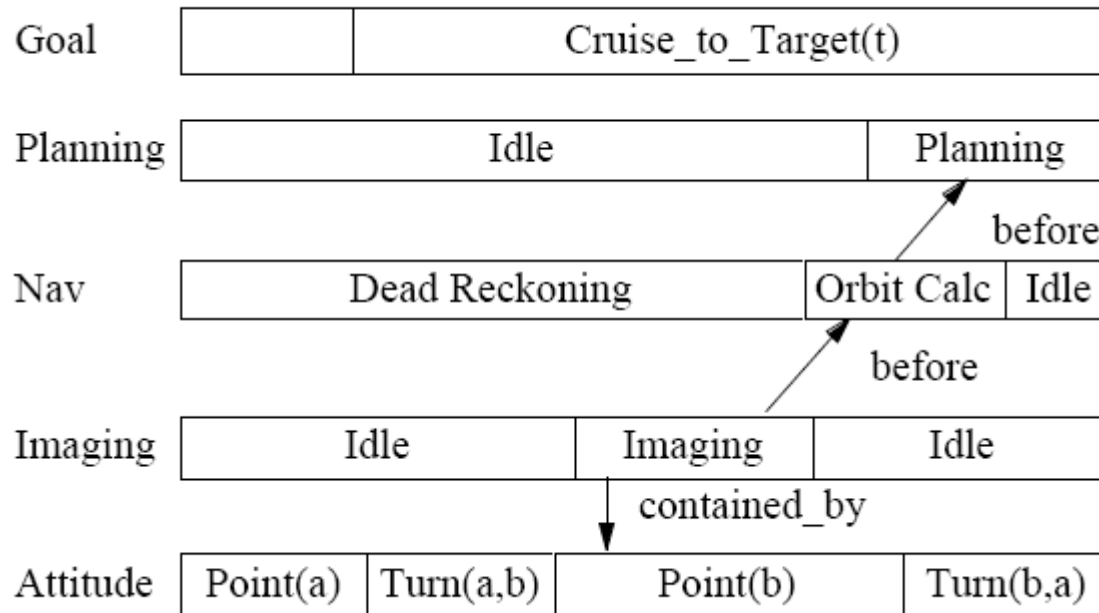
Remote Agent

- Planning



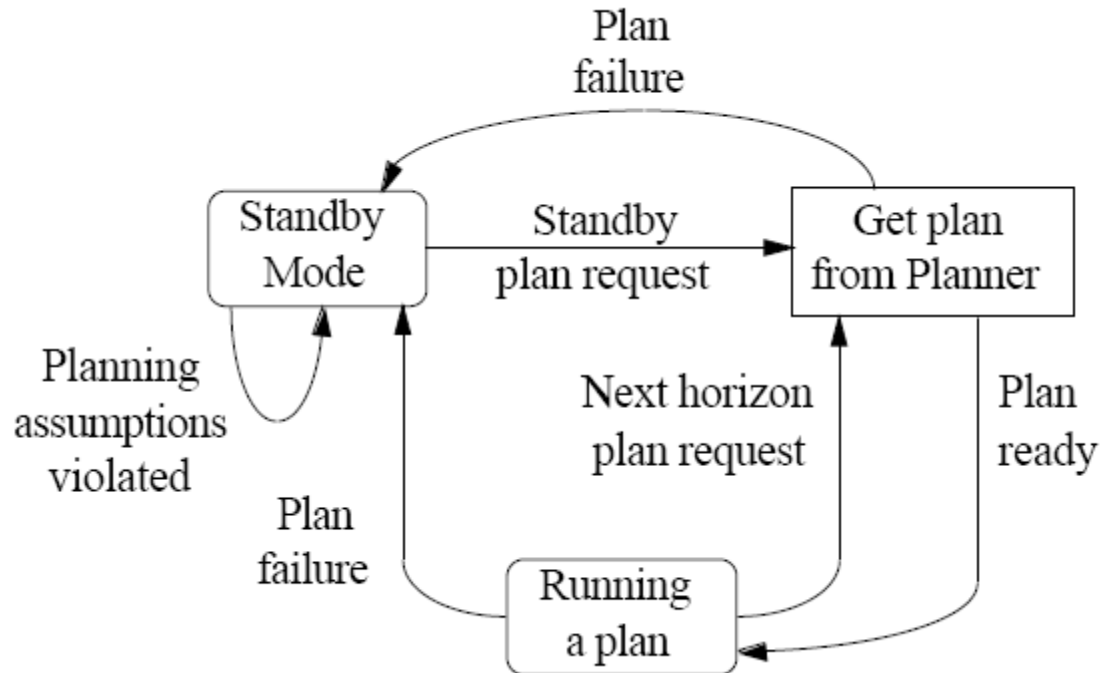
Remote Agent

- Planning to plan



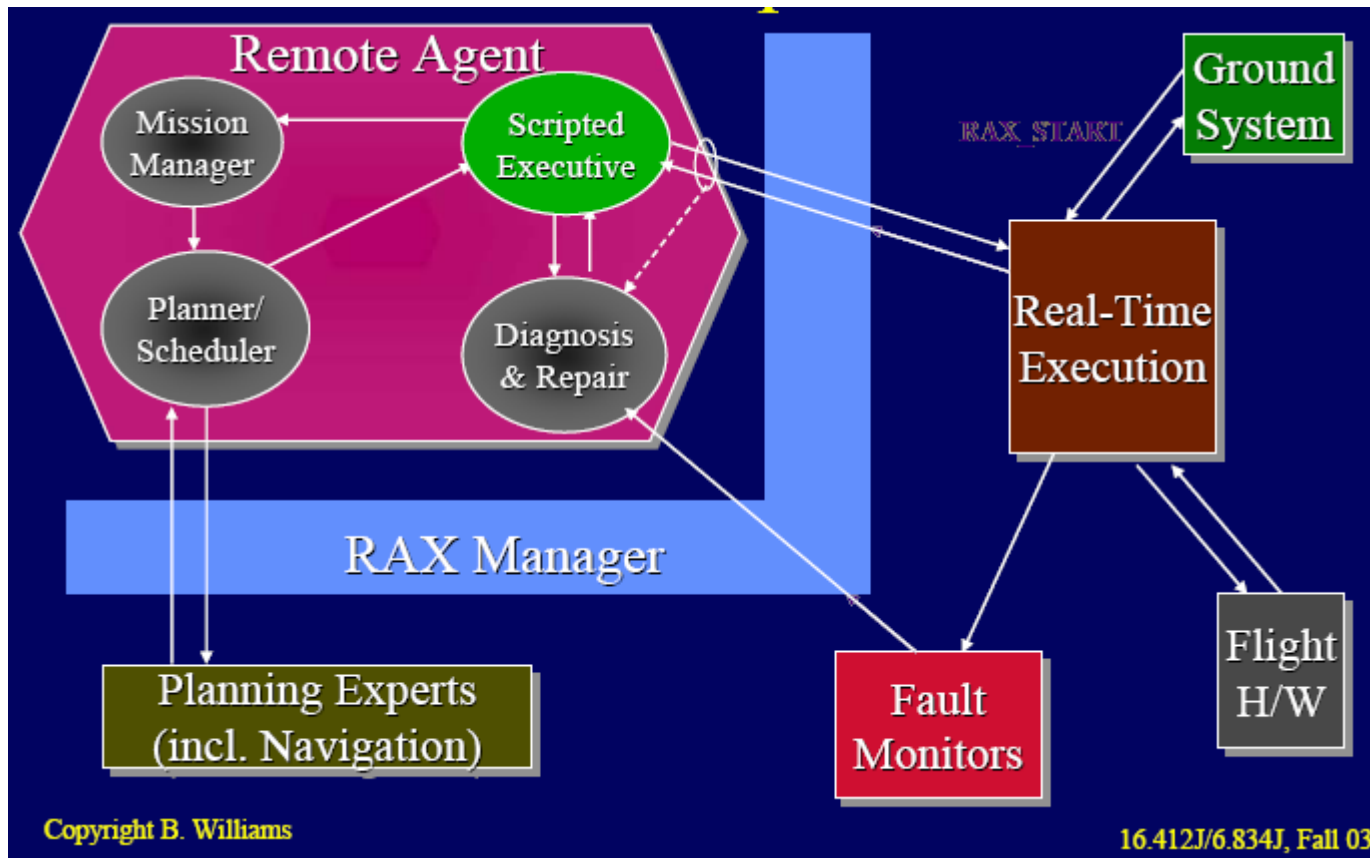
Remote Agent

- Periodic planning and replanning



Remote Agent

- Executive system dispatch tasks

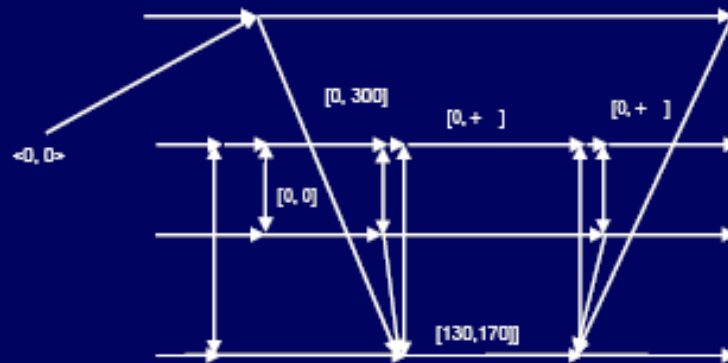


Remote Agent

- The Plan Executor has two duties:
 - Select and Schedule activities for execution
 - Update the network (constraint propagation) after the action execution or execution step (latency)
- Executor Cycle:
 - Activity Graph (STN) from Planner
 - Propagate with latency
 - Enabled time points = scheduled parents (fixed time points)
 - Select and Schedule enabled time points
 - Propagate constraint network given the new binds

Remote Agent

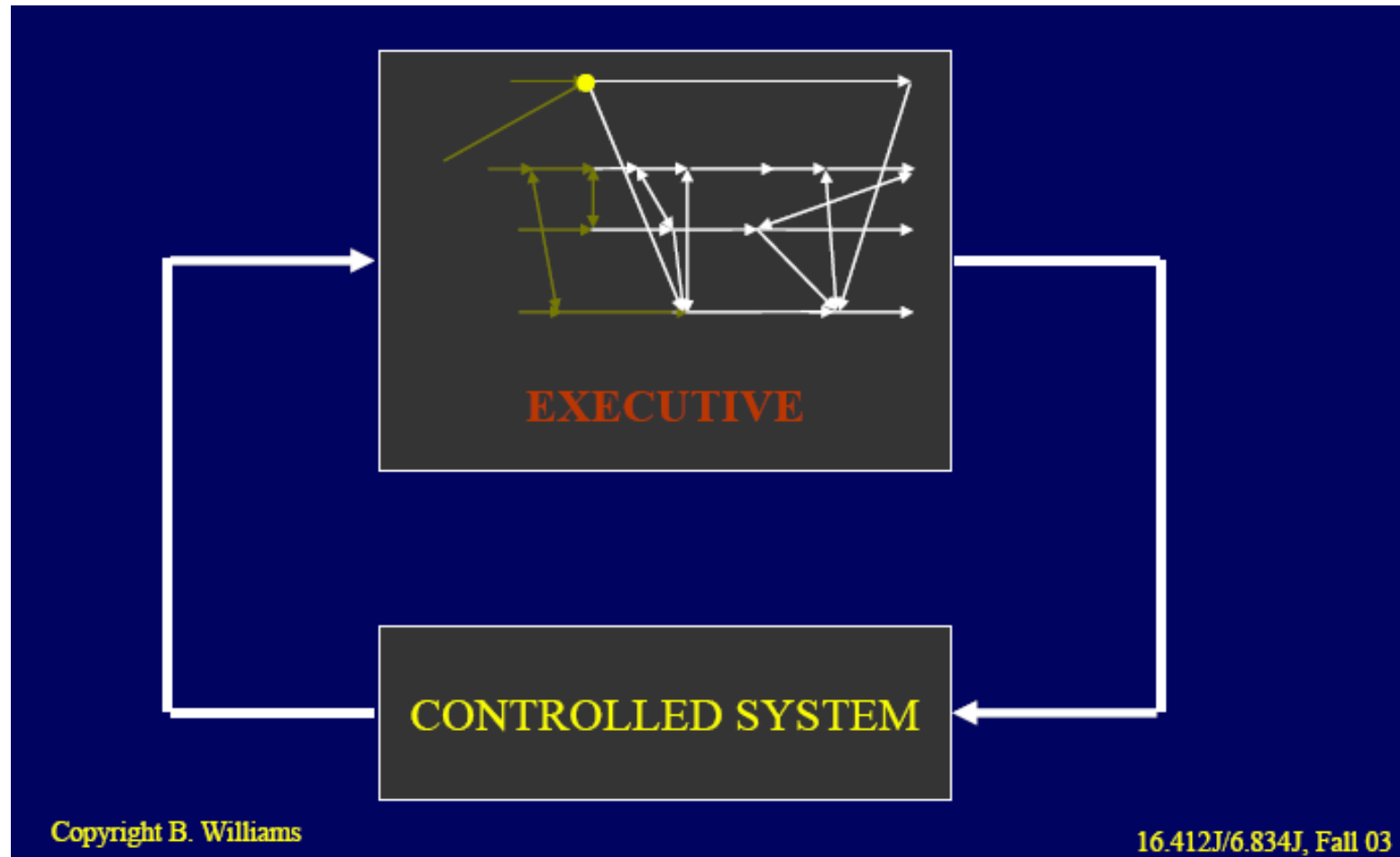
- Executing Flexible Plans



- Propagate temporal constraints
- Select enabled events
- Terminate preceding activities
- Run next activities

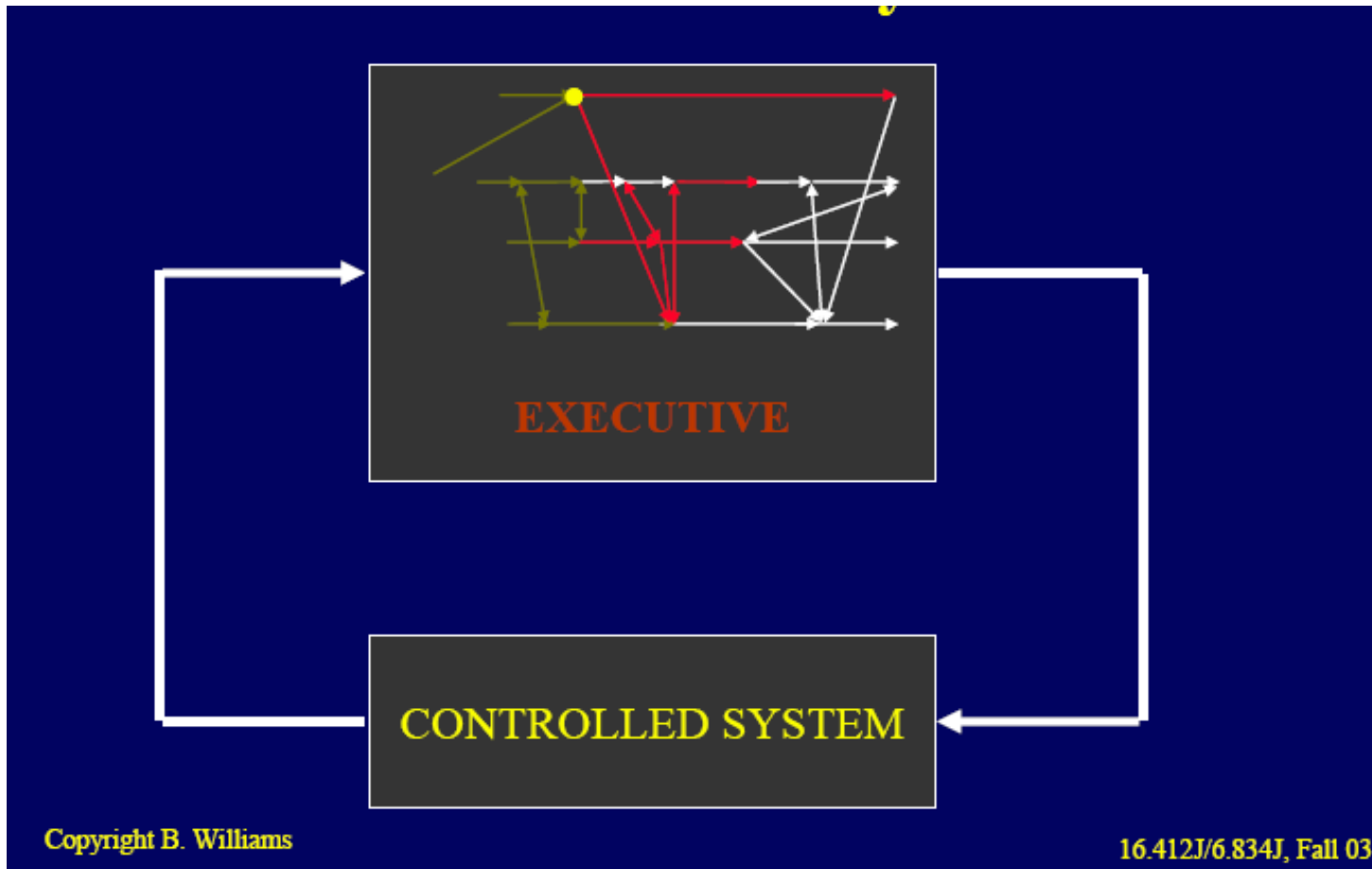
Remote Agent

- Constraint propagation can be costly



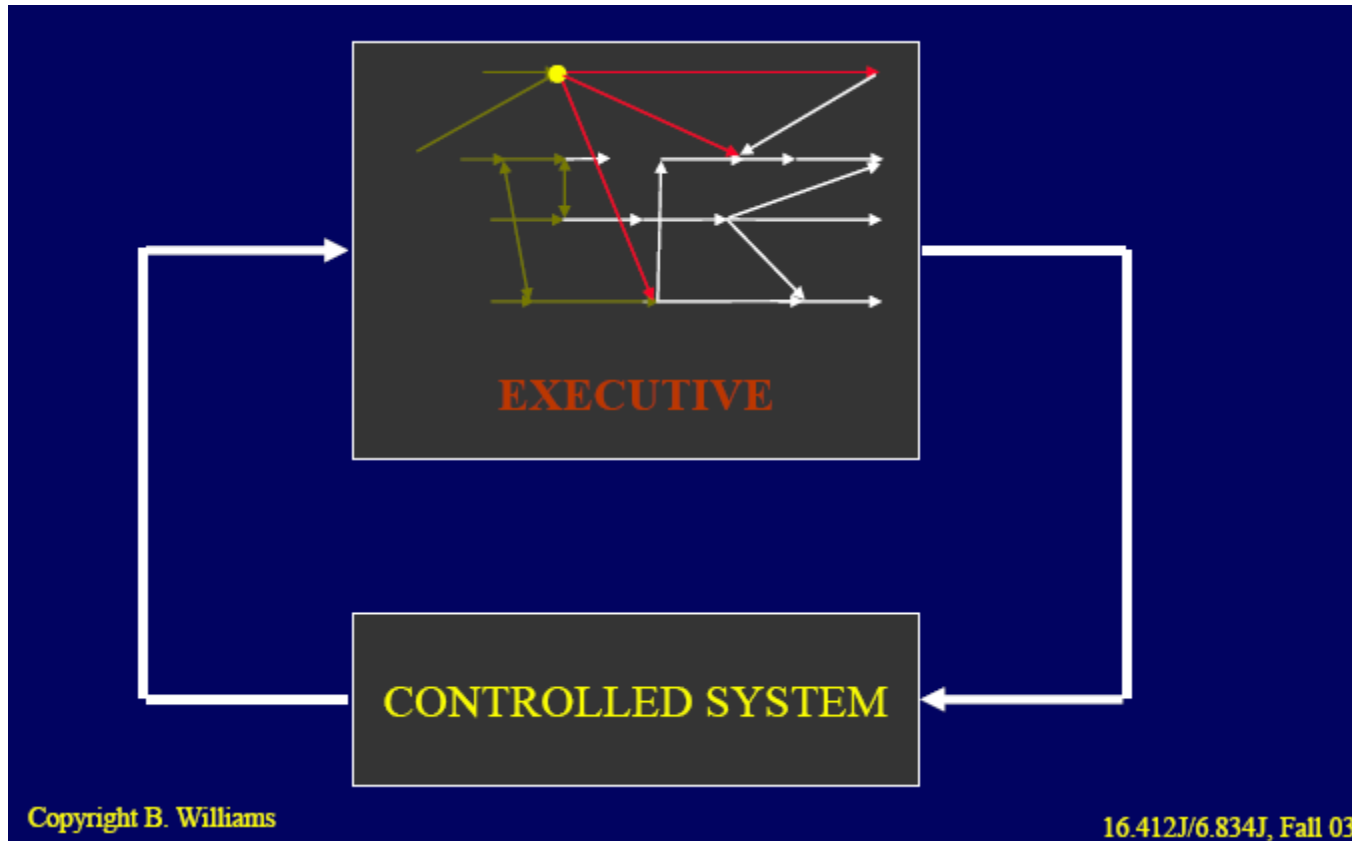
Remote Agent

- Constraint Propagation can be costly



Remote Agent

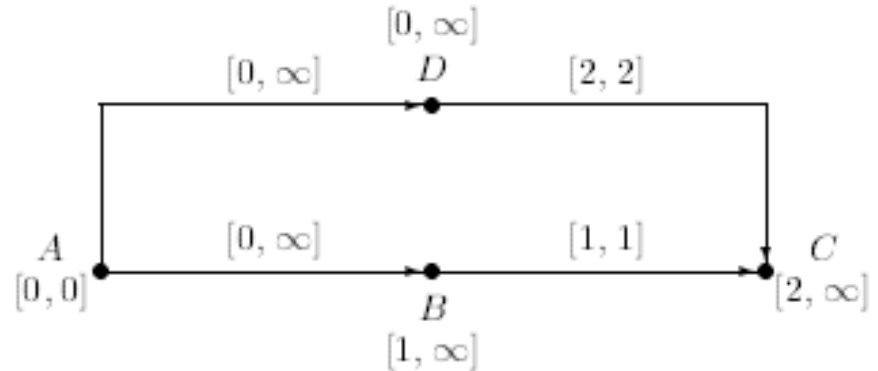
- Solution: compile temporal constraints to an efficient network



Remote Agent

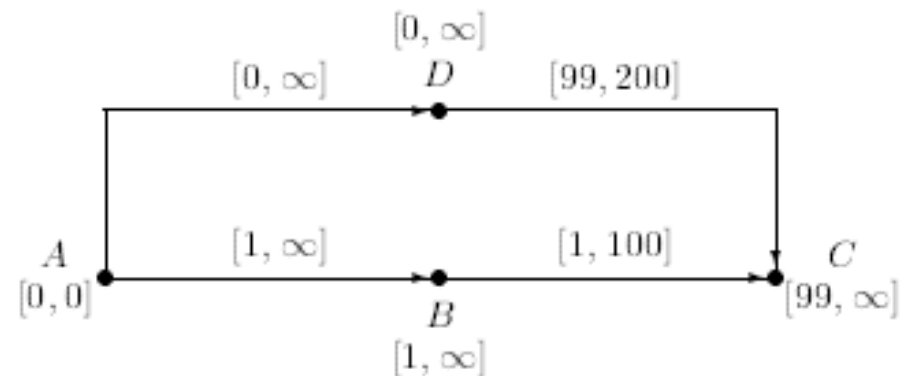
- Dispatchability

- Alcuni vincoli non visibili a tempo di esecuzione;
- Occorre rendere la rete dispatchable aggiungendo vincoli impliciti (e.g. D prima di B)



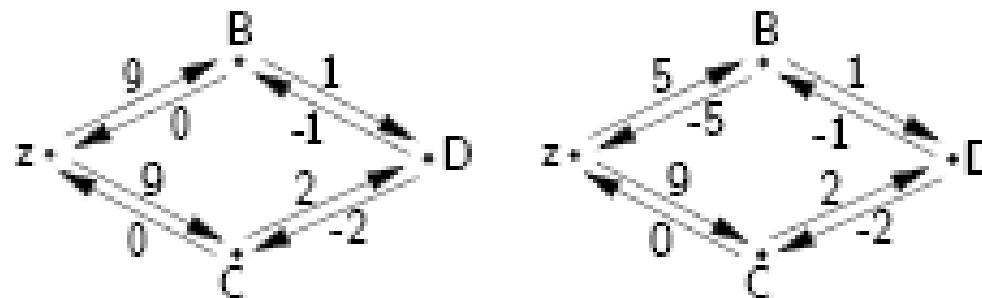
- Compilare la rete in forma dispatchable:

- Introdotti vincoli impliciti
- Tolti vincoli ridondanti



Dispatchability

A Sample Execution*



After executing B at time 5, it turns out that C must be executed at time 4 (which is already past).

* (Muscettola, Morris, & Tsamardinos 1998)

Dispatcher

Greedy Dispatcher*

While some time-points not yet executed:

 Wait until some time-point is executable.

 If more than one, pick one to execute.

 Propagate updates only to *neighboring* time-points (i.e., do not fully update \mathcal{D}).

* (Muscettola, Morris, & Tsamardinos 1998)

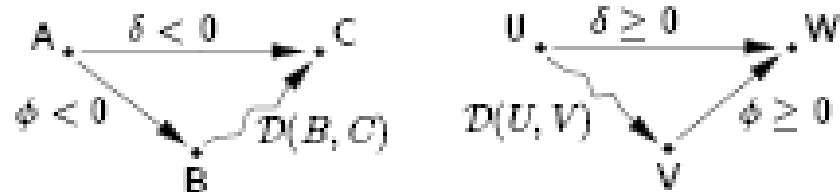
Dispatcher

TIME DISPATCHING ALGORITHM:

1. Let
 - A = {start_time_point}
 - current_time = 0
 - S = {}
2. Arbitrarily pick a time point TP in A such that current_time belongs to TP's time bound;
3. Set TP's execution time to current_time and add TP to S;
4. Propagate the time of execution to its IMMEDIATE NEIGHBORS in the distance graph;
5. Put in A all time points TPx such that all negative edges starting from TPx have a destination that is already in S;
6. Wait until current_time has advanced to some time between
 - min{lower_bound(TP) : TP in A}
 - and
 - min{upper_bound(TP) : TP in A}
7. Go to 2 until every time point is in S.

Dispatchability

Lower and Upper Dominance*



- The *negative edge AC* is *lower-dominated* if:
 $\delta = \phi + \mathcal{D}(B, C)$.
- The *non-negative edge UW* is *upper-dominat'd* if:
 $\delta = \mathcal{D}(U, V) + \phi$.

* (Muscatella, Morris, & Tsamardinos 1998)

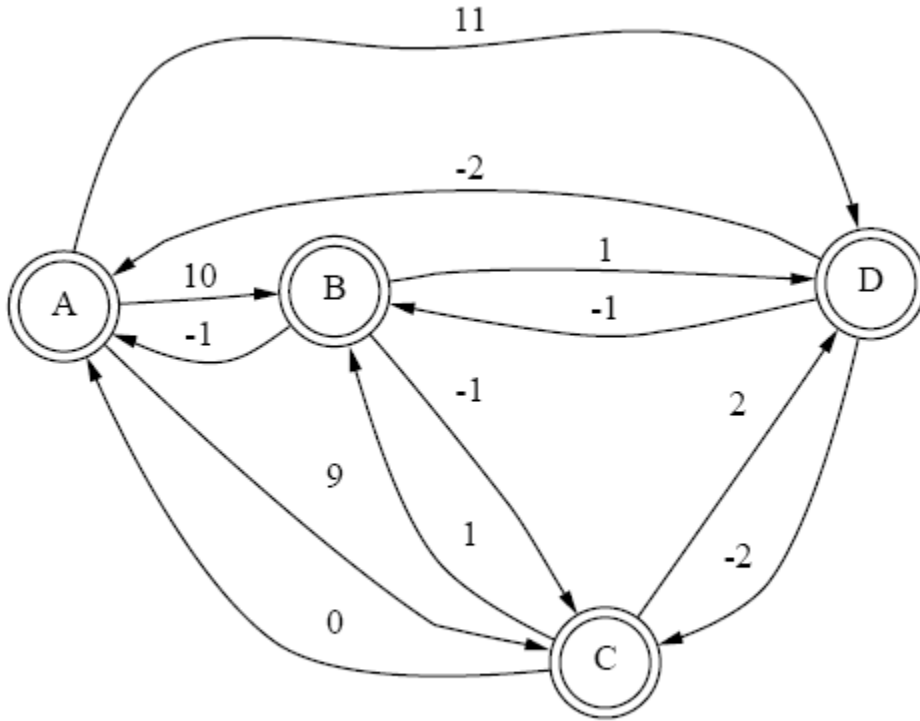
Dispatchability

Dispatchability*

- An STN that is guaranteed to be satisfied by the Greedy Dispatcher is called *dispatchable*.
- Any *consistent* STN can be transformed into an equivalent *dispatchable* STN.
- Step I: The corresponding All-Pairs graph is equivalent and dispatchable.
- Step II: Remove *lower- and upper-dominated* edges (does not affect dispatchability).

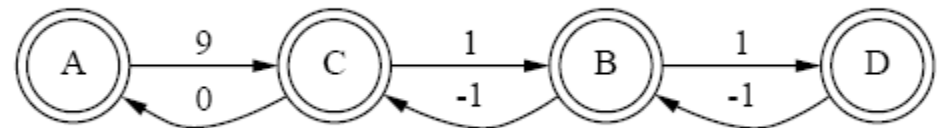
* (Muscettola, Morris, & Tsamardinos 1998).

Dispatchability



All pair graph

Filtered graph



Controllability

- Alcune attività non sono controllabili, ma solo osservabili
- E.g. after start_turn, end_turn ? Quando finisce?
- Il grafo delle attività STN contiene time point controllabili e non controllabili
- Le attività non controllabili non possono essere schedulate, ma solo osservate
- Propagazione??

Controllability

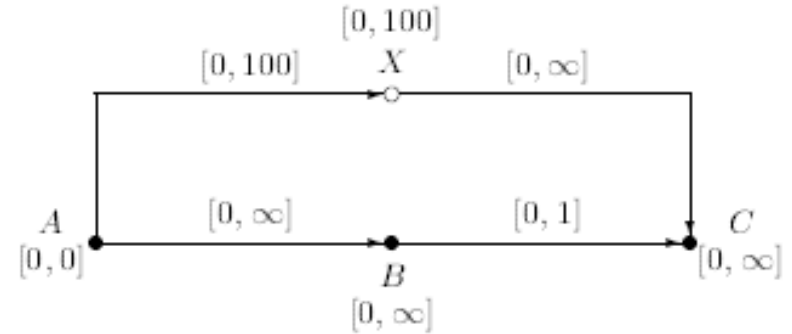
Controllability Issues*

- In real-world applications, an agent may only control some time-points directly; others may be controlled by other agents or Nature.
- Such a network is called *controllable* if there exists a strategy for the agent to execute the time-points under its direct control that will ensure the consistency of the network—no matter how the other agents or Nature execute their time-points.

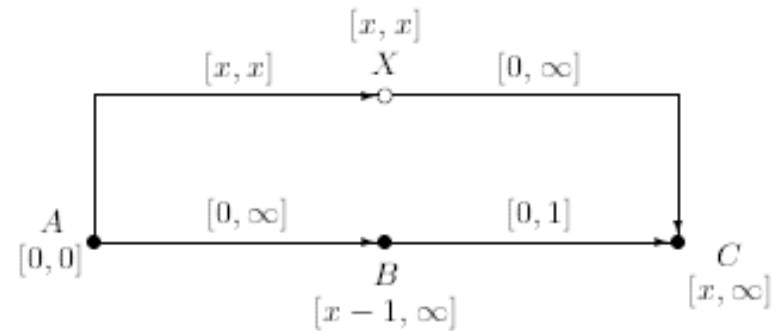
* (Vidal & Ghallab 1995; Vidal & Fargier)

Controllability

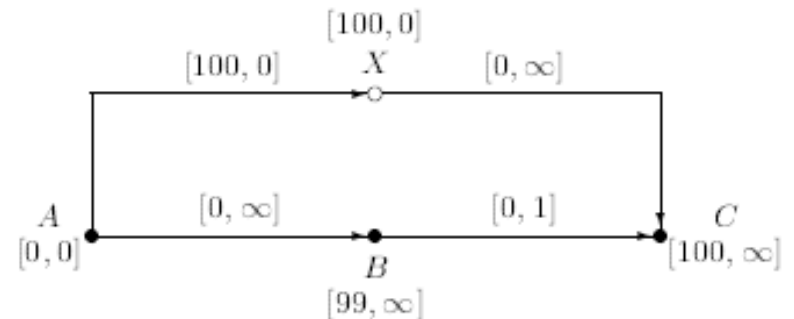
- Gestire eventi non controllabili
- Es. Se B schedulato prima di X, B vincola X



- Soluzione Dinamica:
B dopo X



- Soluzione Forte:
B a 99



Controllability

- **Weak Controllability:** per ogni evento incontrollabile esiste uno scheduling che permette l'esecuzione;
- **Strong Controllability:** esiste uno scheduling robusto qualunque siano gli eventi non controllabili;
- **Dynamic Controllability:** per ogni evento incontrollabile passato esiste uno scheduling che permette l'esecuzione.

Controllability

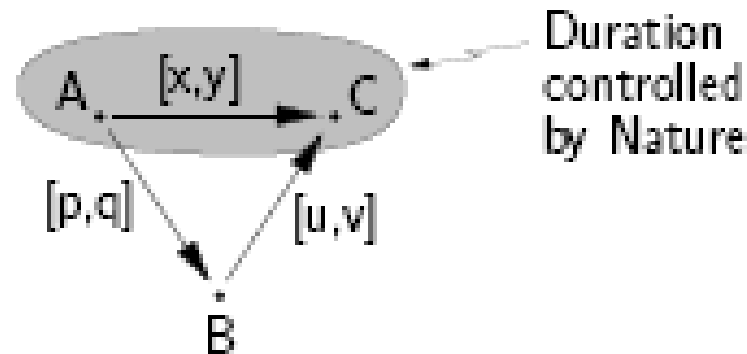
Checking Dynamic Controllability*

Morris et al. (2001) present a sound and complete algorithm for checking dynamic controllability using:

- *Triangular Reductions*
- *Wait Propagation*

Controllability

Triangular Reductions



- If $v < 0$, then B follows C – no reduction.
- If $u \geq 0$, then B precedes C. Must tighten bounds on interval AB to: $[y-v, x-u]$.

Controllability

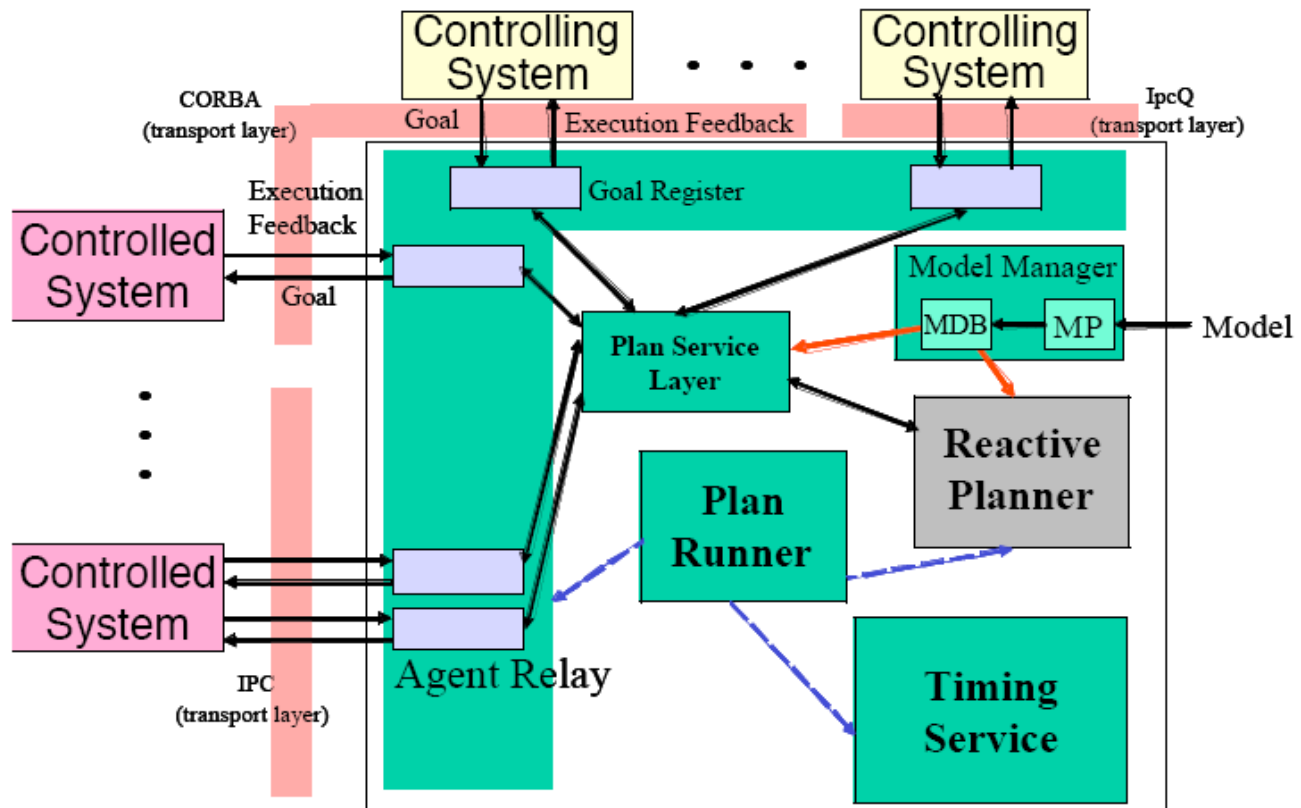
Triangular Reductions (ctd.)

- If $u < 0$ and $v \geq 0$, then the order of B and C is not yet determined. Derive a *WAIT*: If C has not yet been executed, B must wait to be executed until $(y-v)$ after A.

Waits can be propagated much like binary constraints.

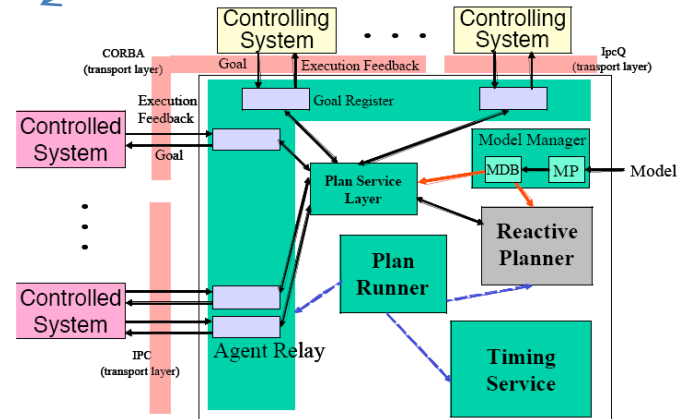
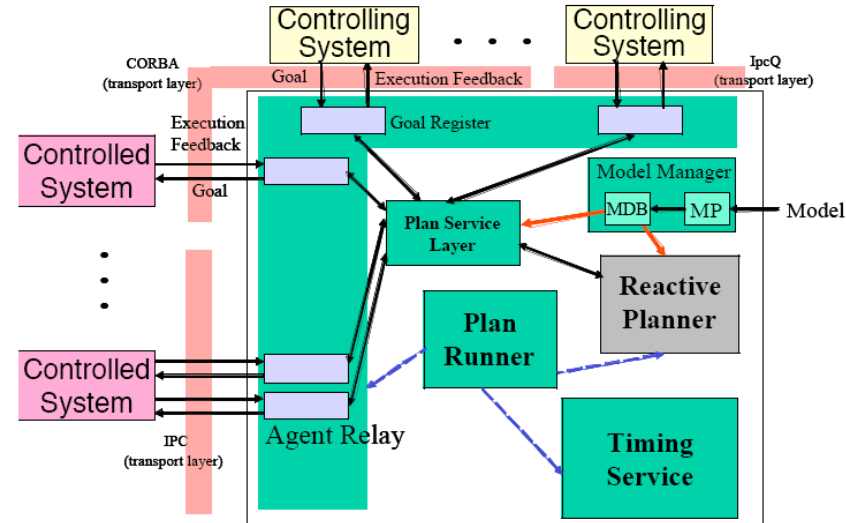
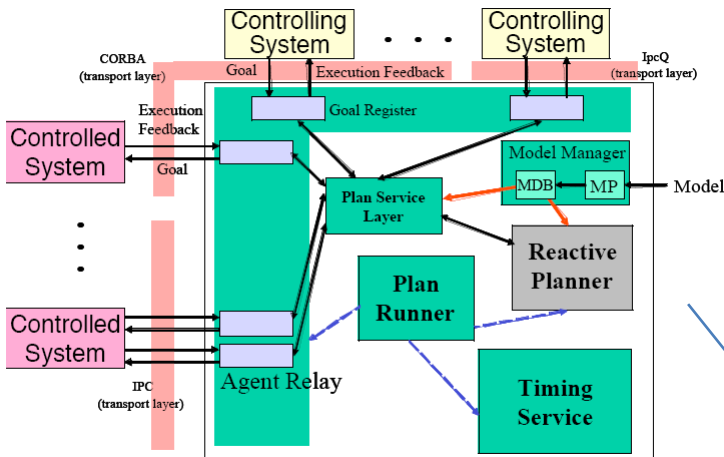
IDEA Architecture

- Evoluzione del RA: reactive and deliberative planning



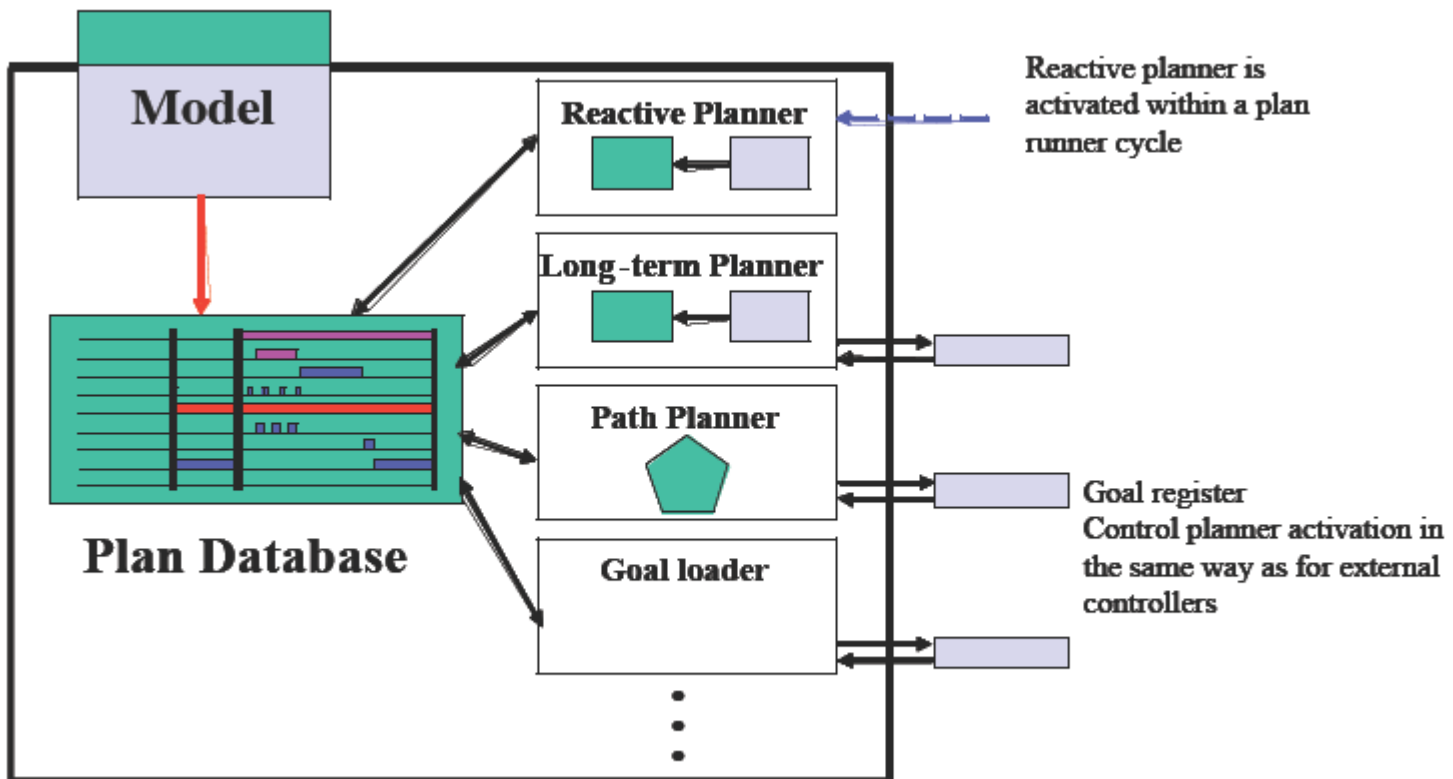
IDEA Architecture

- Multi-agent architecture:



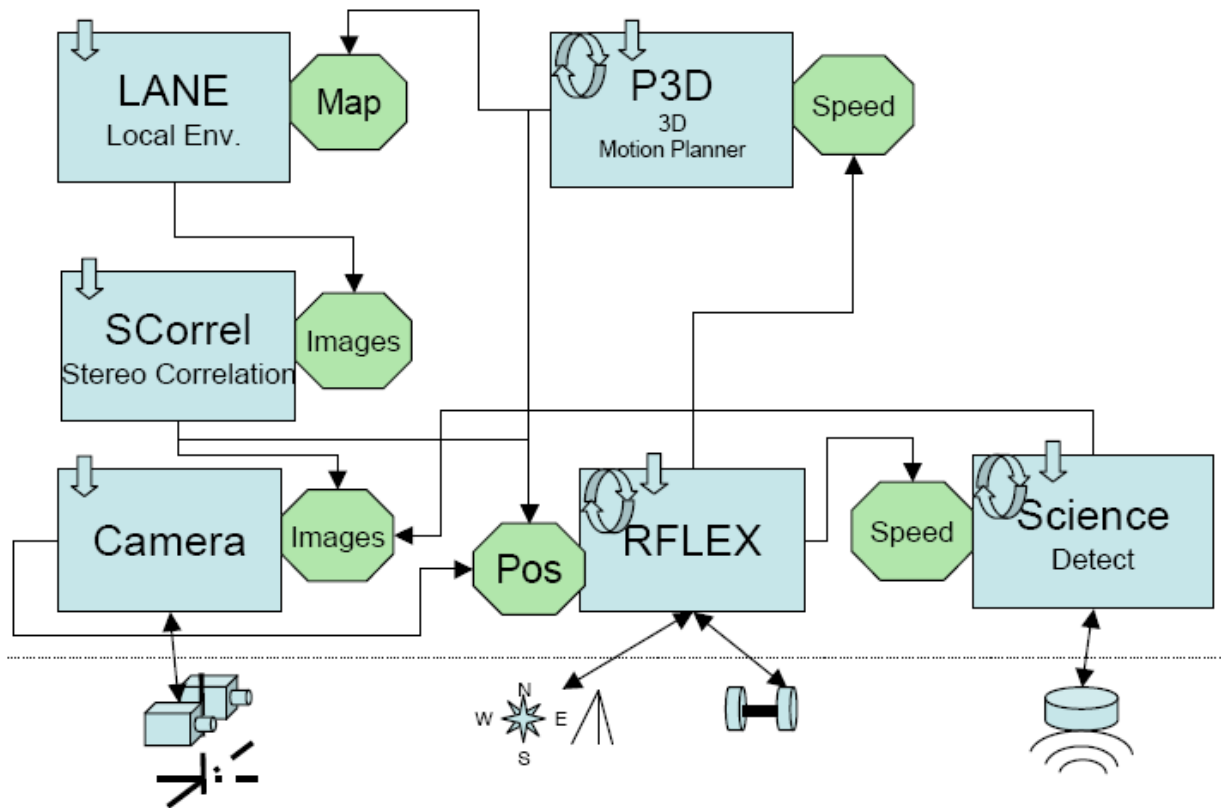
IDEA Architecture

- Reactive planning come controllo
- Interazione deliberative and reactive planning



Functional Layer

- GenoM (LAAS)

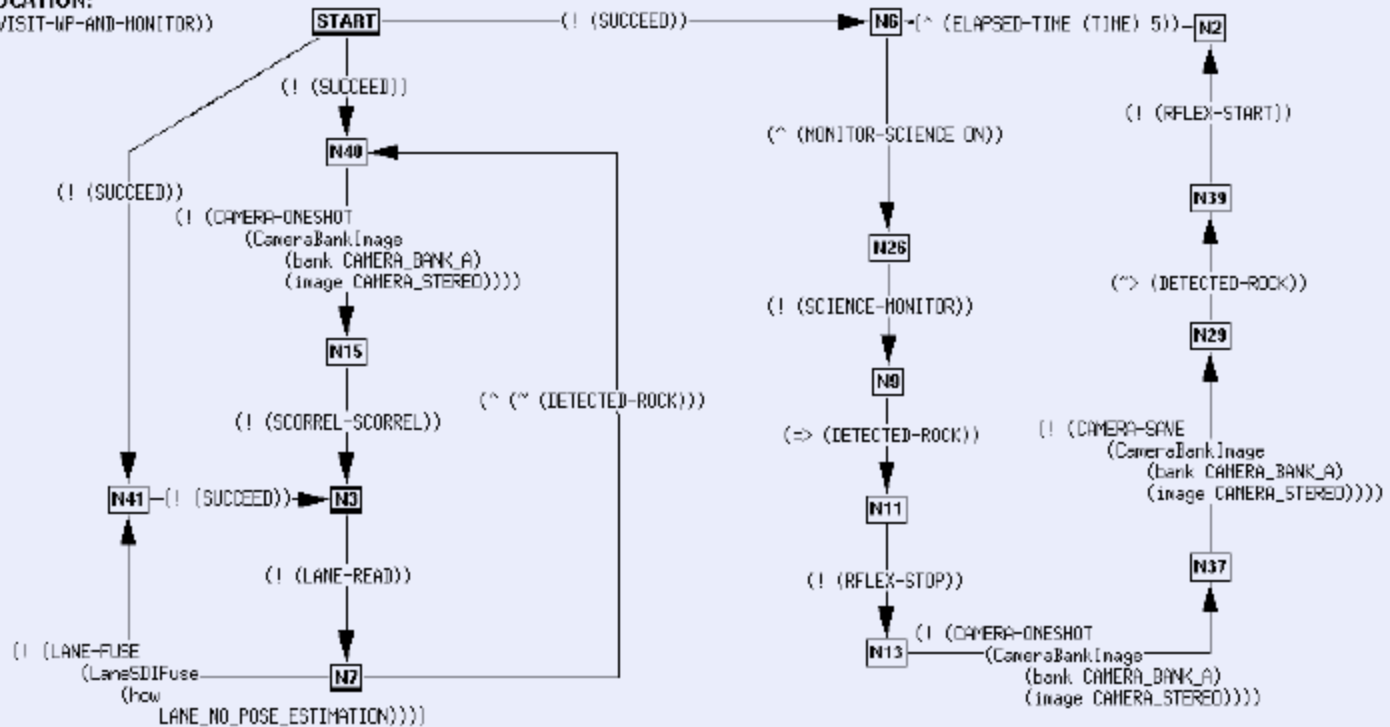


PRS Controller

- A procedural controller (vedi dopo ...)

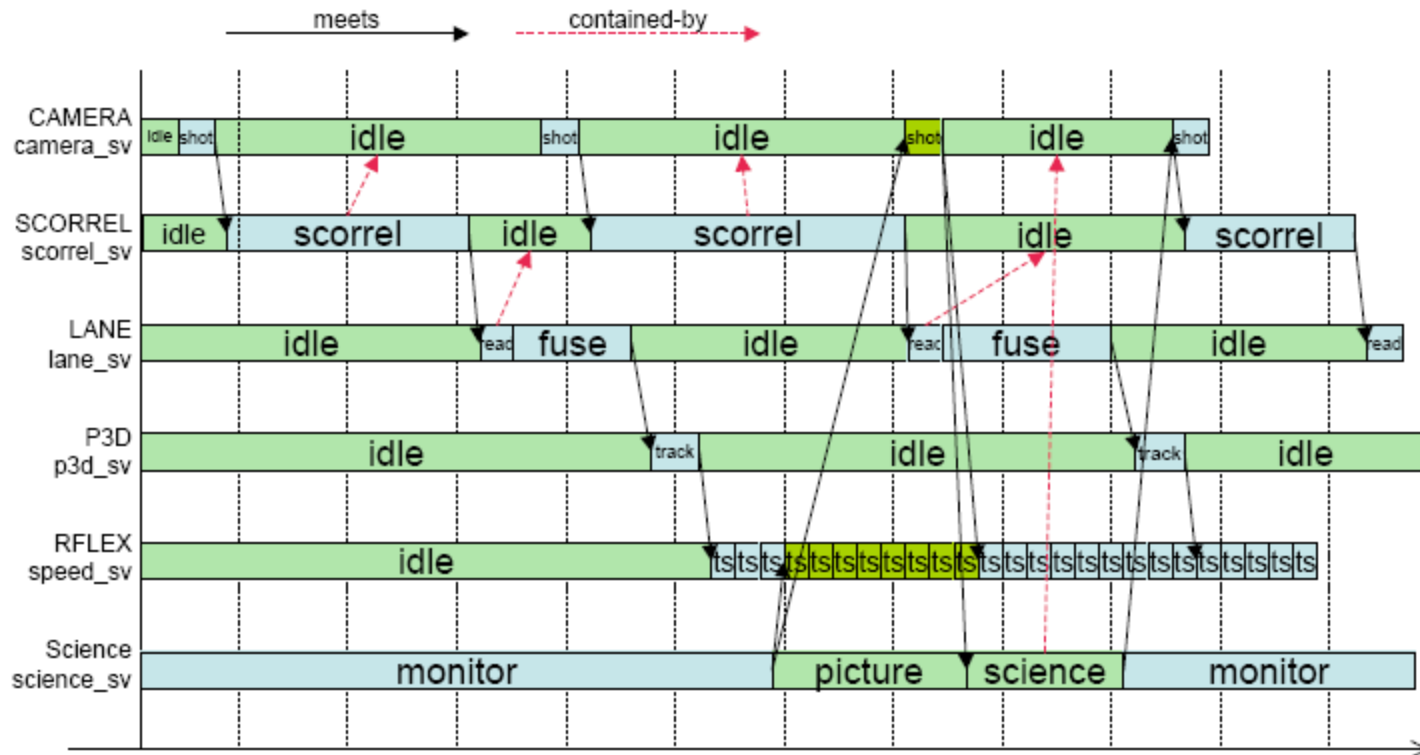
Visit Way Points and Monitor Rocks

INVOCATION:
(! (VISIT-WP-AND-MONITOR))



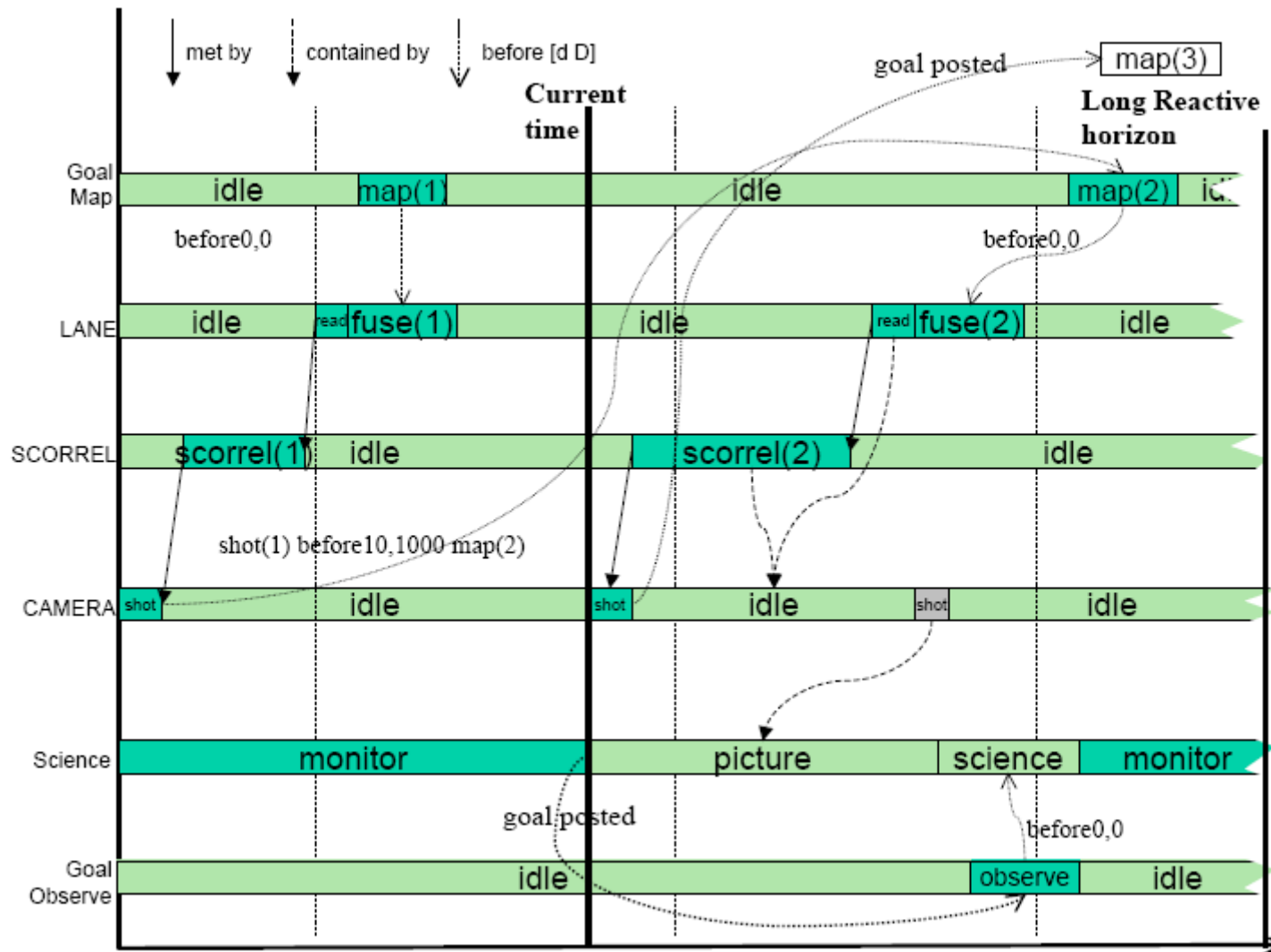
IDEA Architecture

- Attività pianificate (plan database):



IDEA Architecture

- Reactive Planning



IDEA Architecture

- Reactive and Deliberative planning

