

# Fault diagnosis and prognosis in Petri Nets by using a single generalized marking estimation

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# Outline

- 1** Preliminaries
- 2** Motivation
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- 4** Ongoing works

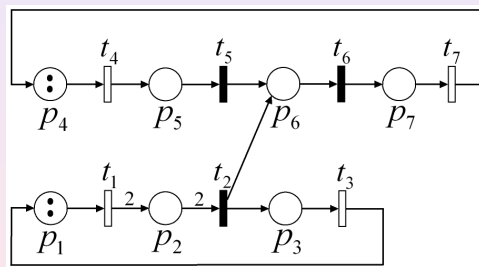
# Backgrounds

- Fault detection for DES has been issued since the mid 80s, and it is still an *hot topic*
- The standard approach is based on the *diagnoser* automata (Sampath et al., IEEE Trans. Aut. Contr., 1995)
- All possible unobservable events that may occur from a given state have to be considered
- A number of approaches based on a Petri net (PN) models have been proposed

## Backgrounds (cont'd)

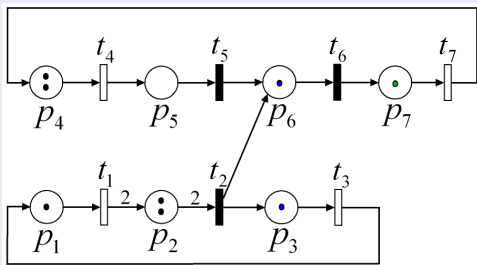
- In the PNs framework, a possible approach to fault diagnosis provides to associate the faults to unobservable transitions
- These approaches need to estimate the current state of the net (Genc and Lafortune, IEEE Trans. Automat. Sci. Eng., 2007 – Giua and Seatzu, 44<sup>th</sup> IEEE CDC, Boel and Jiroveanu, 16<sup>th</sup> Symp. Math Theory Networks Syst.)
- **Explosion of the state space estimation**

# Explosion of the state space estimation



$$\mathbf{m}_0 = [2 \ 0 \ 0 \ 2 \ 0 \ 0 \ 0]^T - t_1 \text{ fires.}$$

# Explosion of the state space estimation



*Unobservable Reach* (as called in Genc and Lafortune)

$$\mathbf{m}_1 = [1 \ 2 \ 0 \ 2 \ 0 \ 0 \ 0]^T$$

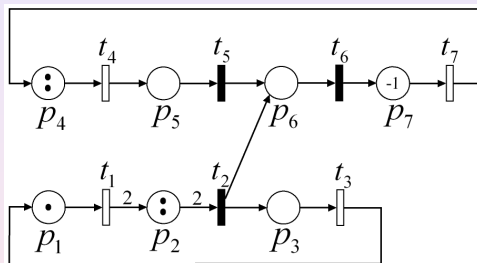
$$\mathbf{m}_2 = [1 \ 0 \ 1 \ 2 \ 0 \ 1 \ 0]^T \text{ - if } t_2 \text{ has fired}$$

$$\mathbf{m}_3 = [1 \ 0 \ 1 \ 2 \ 0 \ 0 \ 1]^T \text{ - if } t_2 \text{ and } t_6 \text{ have fired}$$

# A unique PN state estimation: the generalized marking

- In (Basile et al, WODES 2008) the authors have introduced **generalized markings** to avoid state space explosion.
- **Generalized markings** can have negative components
- The negative components record how many tokens are missing in the input places of observable transitions, whose firings have not been explained yet.
- Using the **generalized marking** the fault diagnosis problem is formulated in terms of ILP problems
- Given the *local* representation of the state in PNs, for each fault the ILPs are solved on a subnet which is *smaller* than the whole plant model.

## Generalized marking: example



If  $t_7$  fires we reach

$$\mu = [1 \ 2 \ 0 \ 2 \ 0 \ 0 \ -1]^T.$$

As far as the fault diagnosis is concerned,  $\mu$  stores in a compact way all the needed information about the state space estimation.



## Fault detection algorithm: remarks

- The problem of diagnosability, i.e. to decide a priori if a given fault can be detected, is not addressed by the proposed algorithm.
- It is assumed that the fault events - assumed to be unobservable - can be detected.
- The proposed approach is mainly aimed to improve the efficiency in terms of memory requirements.

# Motivation

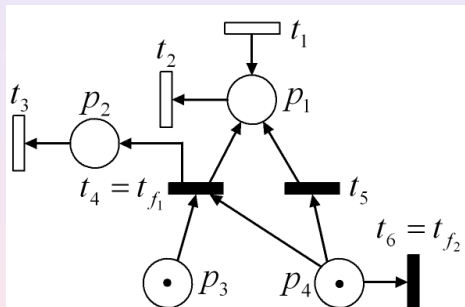
## Motivation

We had the feeling that diagnosability was sufficient to perform *diagnosis* using a **single generalized marking estimation**, but we must prove that!

## Remark

Diagnosability is obviously necessary!

## Motivation (cont'd)



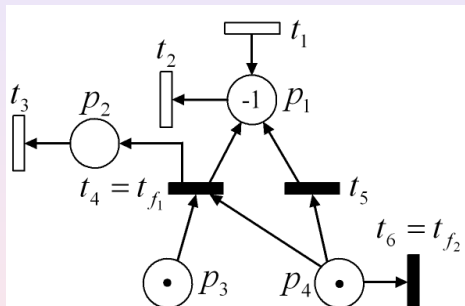
$$\mu_0 = \mathbf{m}_0 = [0 \ 0 \ 1 \ 1]^T$$

$t_1, t_2, t_3$  are observable transitions

$t_4, t_5$  and  $t_6$  are unobservable transitions

$t_4 = t_{f_1}$  and  $t_6 = t_{f_2}$  model faults

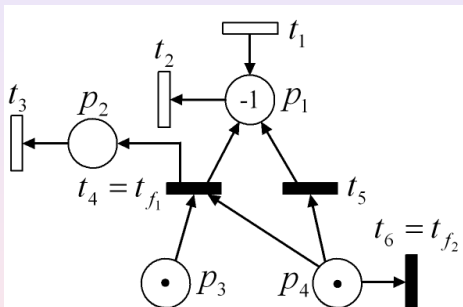
# Motivation



After the firing of  $\sigma = t_2$ , the generalized marking estimation becomes  $\mu_1 = [-1 \ 0 \ 1 \ 1]^T$ .

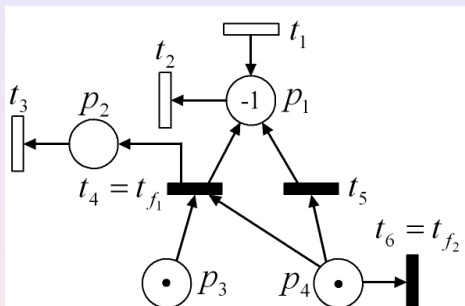
The negative component of  $\mu_1$  means that either  $t_4$  or  $t_5$  should have fired in order to explain the observed firing.

## Motivation - (cont'd)



If  $t_3$  does not fire, it is impossible to find any sufficiently long continuation of  $\sigma$  that permits to distinguish between the firing of  $t_4$  and  $t_5$ , then the language associated with the net is not diagnosable (we will show that it is also not detectable).

## Motivation - (cont'd)

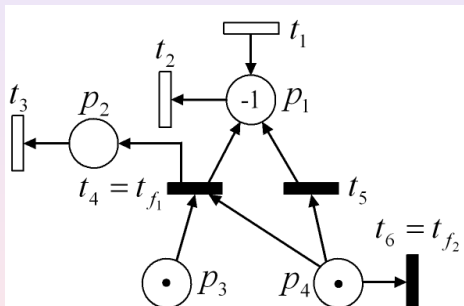


Let  $\Sigma(\mu_1)$  be the set of all the possible firing count vectors  $\epsilon$  corresponding to sequences of unobservable transitions enabled under  $\mu_1$ ,

$$\max_{\epsilon \in \Sigma(\mu_1)} \epsilon(t_{f_2}) = 0,$$

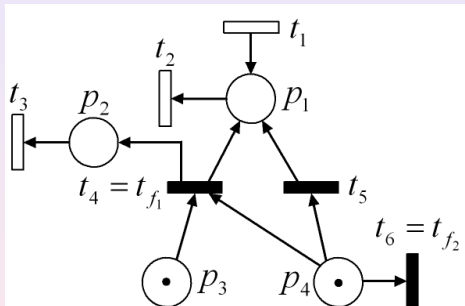
meaning that the fault  $t_{f_2}$  has not occurred for sure (and it cannot occur in the future).

## Motivation - (cont'd)



Moreover  $t_{f_2}$  cannot occur anymore, since either  $t_4$  or  $t_5$  has fired, disabling  $t_{f_2}$  without any possibility to enable it once again.

## Motivation - (cont'd)



If the firing of  $t_1$  is observed  $\mu_2 = \mu_0$  is reached we erroneously get

$$\max_{\epsilon \in \Sigma(\mu_2)} \epsilon(t_{f_2}) = 1,$$

meaning that  $t_{f_2}$  may occur.



# Contribution

We have found the conditions under which a single g-marking estimation can be used to distinguish both

- between “a fault has occurred for sure” and “a fault may not have occurred” (*diagnosis*)
- between “a fault may have occurred/occur” and “a fault has not occurred for sure” (*prognosis*)

## Preliminaries

The notion of detectable prefix-closed and live language is given starting from the definition of diagnosability given in Sampath et al., IEEE Trans. Aut. Contr. 2005.

- $N = (P, T, \mathbf{Pre}, \mathbf{Post})$  is a net with  $T = T_{uo} \cup T_o$ , and  $T_f \subseteq T_{uo}$ .
- $\bar{s}$  is the prefix-closure of any trace  $s \in T^*$ . We denote by  $L/s$  the post-language of  $L$  after  $s$ .
- $Pr : T^* \mapsto T_o^*$  is the usual projection which “erases” the unobservable events in a trace  $s$ .
- $Pr_L^{-1}$  is the inverse projection operator defined as

$$Pr_L^{-1}(r) = \{s \in L \text{ s.t. } Pr(s) = r\}.$$

- If  $\dot{t}$  is the final event of trace  $s$ , we define

$$\Psi(t_{f_i}) = \{st \in L \text{ s.t. } \dot{t} = t_{f_i}\}.$$

## Diagnosable language - Definition

A prefix-closed and live language  $L$  is said to be diagnosable w.r.t.  $T_f$  if

$\forall t_{f_i} \exists h_i \in \mathbb{N}$  such that the following holds

$\forall s \in \Psi(t_{f_i})$  and  $\forall q \in L/s$

$\|q\| \geq h_i \Rightarrow D$

where  $\|q\|$  is the length of trace  $q$ , and the diagnosability condition  $D$  is

$r \in Pr_L^{-1}(Pr(sq)) \Rightarrow t_{f_i} \in r$ .

Let  $s$  be any trace generated by the system that ends in a failure event  $t_{f_i}$ , and let  $q$  be any sufficiently long continuation of  $s$ . Condition  $D$  implies that along every continuation  $q$  of  $s$  it is possible to detect the occurrence of  $t_{f_i}$  with a finite delay, specifically in at most  $h_i$  transitions of the system after  $s$ .

# Detectable language

## Definition

A prefix-closed and live language  $L$  is said to be detectable w.r.t.  $T_{uo}$  if it is diagnosable w.r.t.  $T_{uo}$ .

## Remarks

- detectability implies diagnosability
- undetectability does not necessarily implies undiagnosability
- undiagnosability implies undetectability

# Main result - 1

## Theorem

Let  $L$  be diagnosable w.r.t.  $T_f$ . If  $s \in \Psi(t_{f_i})$  then exists  $q \in L/s$ , such that

$$\min_{\epsilon \in \Sigma(N, \mu)} \epsilon(t_{f_i}) > 0,$$

with  $\mu_0[z] \mu$ , and  $z = Pr(sq)$ .

## Main result - 2

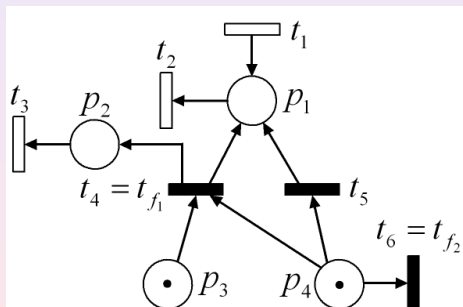
### Theorem

Let  $L$  be detectable w.r.t.  $T_{uo}$ . If  $s$  is a sequence which enables the firing of  $t_{f_i}$  and  $t_{f_i} \notin s$ , then it exists  $h \in \mathbb{N}$  such that for all sequences  $q \in L/s$  whose firing does not enable  $t_{f_i}$ , and  $\|q\| > h$ ,  $t_{f_i} \notin q$ , it holds that

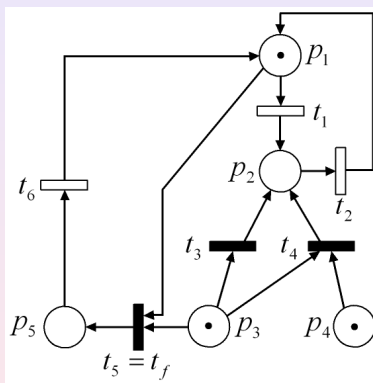
$$\max_{\epsilon \in \Sigma(N, \mu'')} \epsilon(t_{f_i}) = 0,$$

with  $\mu_0 [\beta] \mu''$ , and  $\beta = Pr(sq)$ .

# Undiagnosable and undetectable net



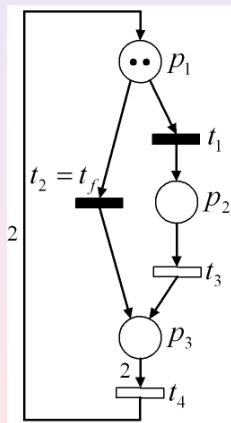
# Diagnosable and undetectable net



- It is still not possible to distinguish between the firing of  $t_3$  and  $t_4$ , hence the language is undetectable.
- The language is diagnosable. Indeed after the firing of  $t_f$  all the possible continuations are given by  $\dots t_f t_6 (t_1 t_2)^*$ .



# Detectable net



## Ongoing works

- An updated version of the fault detection algorithm based on **g-markings**, which includes the present results, has been published in [Basile et al., IEEE Trans. Aut. Contr., Apr. 2009](#)
- We have proposed a new approach for fault diagnosis based on ILPs without using the g-markings (see [Basile et al., IFAC DCSD 2009, Jun. 2009](#)). **In this case the detectability assumption is no more needed**
- We are now working on the identification issue to face the problem of fault diagnosis when the fault are not modeled (see [Basile et al., 14<sup>th</sup> IEEE ETFA, Sep. 2009](#))

... **The End**

Thank you!